

FORM PTO-1390  
(REV. 5-93)U.S. DEPARTMENT OF COMMERCE  
PATENT AND TRADEMARK OFFICE

ATTORNEY'S DOCKET NUMBER

H-996

TRANSMITTAL LETTER TO THE UNITED STATES  
DESIGNATED/ELECTED OFFICE (DO/EO/US)

U.S. APPLICATION NO. (if known, see 37 CFR 1.51)

09/889410

International Application No.

PCT/JP00/01333

International Filing Date

March 6, 2000

Priority Date Claimed

March 19, 1999

Title of Invention INFORMATION PROCESSING APPARATUS

Applicant(s) for DO/EO/US J. KITAHARA et al (see attached)

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a FIRST submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a SECOND or SUBSEQUENT submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
  - a. ☐ is transmitted herewith (required only if not transmitted by the International Bureau).
  - b. ☒ has been transmitted by the International Bureau.
  - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US).
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☐ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
  - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
  - b. ☐ have been transmitted by the International Bureau.
  - c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
  - d. ☐ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☒ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern other document(s) or information included:

11. ☒ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☐ A FIRST preliminary amendment.  
☐ A SECOND or SUBSEQUENT preliminary amendment.
14. ☐ A substitute specification.
15. ☐ A change of power of attorney and/or address letter.
16. ☒ Other items or information:

☒ LIST OF INVENTORS' NAMES AND ADDRESSES.

U.S. Application No. (if known, see 37 CFR 1.5) <b>09/889410</b>		International Application No. PCT/JP00/01333		Attorney's Docket Number H-996	
---	--	---	--	-----------------------------------	--

17. <input checked="" type="checkbox"/> The following fees are submitted:				CALCULATIONS	PTO USE ONLY
<u>Basic National Fee (37 CFR 1.492 (a)(1)-(5)):</u>  Search Report has been prepared by the EPO or JPO ..... \$860.00 International preliminary examination fee paid to USPTO (37 CFR 1.482) ..... \$690.00 No international preliminary examination fee (37 CFR 1.482) but international search fee paid to USPTO (37 CFR 1.445 (A)(2)) ..... \$710.00 Neither international examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(A)(2)) paid to USPTO ..... \$1000.00  International preliminary examination fee paid to USPTO (37 CFR 1.482) and all claims satisfied provisions of PCT Article 33(2) to (4) ..... \$100.00					
<b>ENTER APPROPRIATE BASIC FEE AMOUNT = \$ 860.00</b>					
Surcharge of \$130.00 for furnishing the oath or declaration later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(e)).				+ \$ 0.00	
Claims	Number Filed	Number Extra	Rate		
Total	12 -20 =	0	x \$18.00	\$ 0.00	
Independent	2 - 3 =	0	x \$80.00	\$ 0.00	
Multiple dependent claim(s) (if applicable)				+ \$270.00 \$ 0.00	
<b>TOTAL OF ABOVE CALCULATIONS =</b>				<b>\$ 860.00</b>	
Reduction by 1/2 for filing by small entity, if applicable. Verified Small Entity statement must also be filed. (Note 37 CFR 1.9, 1.27, 1.28).				\$ 0.00	
<b>SUBTOTAL =</b>				<b>\$ 860.00</b>	
Processing fee of \$130.00 for furnishing the English translation later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(f)).				+ \$ 0.00	
<b>TOTAL NATIONAL FEE =</b>				<b>\$ 860.00</b>	
Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). \$40.00 per property.				+ \$ 0.00	
<b>TOTAL FEES ENCLOSED =</b>				<b>\$ 860.00</b>	
				Amount to be:	
				Refunded \$	
				Charged \$	

a. ☒ A check in the amount of \$ 860.00 to cover the above fees is enclosed.

b. ☐ Please charge my Deposit Account No. 50-1417 in the amount of \$ \_\_\_\_\_ to cover the above fees.  
A duplicate copy of this sheet is enclosed.

c. ☒ The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any  
overpayment to Deposit Account No. 50-1417. A duplicate copy of this sheet is enclosed.

**Note:** Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive  
(37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status.


SEND ALL CORRESPONDENCE TO:

MATTINGLY, STANGER & MALUR, P.C.

1800 Diagonal Rd., Suite 370

Alexandria, Virginia 22314

(703) 684 -1120

  
 Signature  
 John R. Mattingly  
 Name  
  
 30,293  
 Registration Number

23/PRTS

JG17 Rec'd PCT/PTO 17 JUL 2001

09/889410

-1-

## SPECIFICATION

### INFORMATION PROCESSING APPARATUS

#### TECHNICAL FIELD

The present invention relates to an information processing apparatus using cryptography for ensuring the security of sensitive information in storage and transfer thereof, and more particularly to an information processing architecture for providing high-level information security protection.

#### BACKGROUND ART

The following conventional techniques for an information processing apparatus using cryptography are known.

In JP-A-275115/1998, there is disclosed a cryptographic technique of encrypting information and storing the encrypted information into an external storage device such as a hard disk drive. In this conventional cryptographic technique, when encrypted data Ya and Yb stored in an external storage device 12 are transferred to an information terminal unit 11, the encrypted data Ya and Yb are decrypted in succession using a deciphering key Kb held in an enciphering/deciphering key custody area 35.

09889410-071701

JP-A-214233/1998 discloses another cryptographic technique of providing a dedicated cryptographic processor in an information processing apparatus. According to this conventional cryptographic technique, the cryptographic processor for encrypting data to generate a body of a secure file is equipped in a portable-type personal computer.

In common applications of encryption and decryption processing, sensitive information to be kept in secrecy resides in main memory since data processing is carried out on the main memory. For information encryption in which information is processed through an encryption algorithm, it is required to securely handle the encryption algorithm, cipher key information, and sensitive information to be encrypted.

The above-cited conventional cryptographic techniques, however, give rise to the following problems:

In the conventional cryptographic techniques, since sensitive information and intermediate results of encryption processing are held in main memory, there is a problem that the sensitive information may be extracted by any means. In an information processing apparatus comprising a plurality of semiconductor devices serving as a CPU, main memory and other circuits, when encryption processing is carried out under control of the CPU, an encryption algorithm, sensitive information to be

05889410-071701

Further, since the information processing apparatus contains signal lines (e.g., bus lines) for interconnecting semiconductor device components thereof, data prior to encryption or decrypted data may be easily extracted by means of tapping the signal line and analyzing data attained through the tapping.

## DISCLOSURE OF THE INVENTION

Encryption processing is carried out within a semiconductor device contained in the information processing apparatus. Key information necessary for

encryption is also generated in the semiconductor device. Information regarding encryption is not output onto signal lines arranged in the information processing apparatus. Information which does not affect security protection even if tapped by any other person, such as encrypted information or information not requiring encryption, is output onto the signal lines in the information processing apparatus. The information regarding encryption includes information which is not yet encrypted and information to be used for decryption.

More specifically, according to one aspect of the present invention, there is provided an information processing apparatus in which a semiconductor chip serving as a CPU thereof integrally contains a RAM, a cryptographic processing algorithm memory, a cryptographic processing hardware circuit, a key information generating hardware circuit, and a key information storage hardware circuit. In the present invention, the term "CPU" is used for the sake of convenience. Any semiconductor chip in the information processing apparatus is applicable in lieu of the CPU. In particular, it is preferable to adopt a processor which carries out control and computational operations in the information processing apparatus. That is to say, the present invention provides an information processing apparatus comprising a semiconductor chip in which

05885410-071701

encryption processing including generation of key information is performed in a closed fashion. Further, according to another aspect of the present invention, there may be provided such an arrangement that encryption processing is performed in each of a plurality of CPUs.

Furthermore, according to another aspect of the present invention, encryption processing may be performed in the RAM contained in the CPU.

Moreover, according to another aspect of the present invention, the RAM contained in the CPU may be used as main memory for execution of application programs as well as encryption processing.

Still further, according to another aspect of the present invention, there may be provided such an arrangement that application programs themselves are encrypted and an external storage device holds encrypted files thereof.

Still further, according to another aspect of the present invention, an external bus controller may be provided for controlling data output to an external bus. The external bus controller may control data output so that data is not output to the external bus while the RAM contained in the CPU is accessed. Further, the external bus controller may form a judgment on whether or not to output data to the external bus. If it is judged that data is to be output, the data is delivered to the external bus.

09889410-074704

Still further, according to another aspect of the present invention, encryption and decryption of communication data may be performed in the CPU.

Still further, in any of these arrangements mentioned above, it is also preferable to provide means for determining whether or not to encrypt data according to the contents thereof. If encryption is not necessary for data, the data is output onto the signal lines in the information processing apparatus.

Besides, in a modified arrangement of the present invention, encryption processing may be performed in a processor of a disk system controller for encrypting file location information on a magnetic disk.

#### BRIEF DESCRIPTION OF DRAWINGS

FIG. 1 is a diagram showing a configuration of an information processing apparatus according to the present invention. FIG. 2 is an explanatory diagram of file generation in the information processing apparatus of the present invention. FIG. 3 is a diagram showing an information processing apparatus configuration in which a CPU contains main memory according to a preferred embodiment of the present invention. FIG. 4 is a diagram showing an information processing apparatus configuration in which an application program held in an external storage device is

09889410-074701



encrypted by a CPU according to a preferred embodiment of the present invention. FIG. 5 is a diagram showing a configuration of an external bus controller. FIG. 6 is an explanatory diagram of a preferred embodiment of the present invention in which data output to an external bus is prevented by the external bus controller. FIG. 7 is a diagram showing a configuration of a random number generator required for key generation. FIG. 8 is a diagram showing a configuration of a key custody area. FIG. 9 is a diagram showing a relationship between encryption/decryption processing and key information in a case where encryption and decryption are carried out in the same apparatus. FIG. 10 is a diagram showing a configuration in which the amount of key information to be stored is reduced through the use of a different key handling procedure. FIG. 11 is a diagram showing a relationship between encryption processing and key information and a relationship between decryption processing and key information in a case where different apparatuses are used for encryption and decryption respectively. FIG. 12 is a diagram showing a configuration in which authentication information of a sender is added to the arrangement illustrated in FIG. 11. FIG. 13 is a diagram showing a mechanism for authenticating key information attained from a communicating party. FIG. 14 is a diagram showing a configuration in a case where the present

invention is applied to a processor bus and a system information processing apparatus. FIG. 15 is a diagram showing a configuration in a case where the present invention is employed for a communication application. FIG. 16 is an explanatory diagram showing a configuration in a case where the present invention is applied to an external storage device. FIG. 17 is a diagram for explaining a flow of writing encrypted file location information in the configuration illustrated in FIG. 16. FIG. 18 is a diagram showing a configuration of a disk controller. FIG. 19 is a diagram showing an information processing apparatus comprising a plurality of CPUs in a preferred embodiment of the present invention. FIG. 20 is a diagram showing a modified embodiment of the arrangement illustrated in FIG. 19. FIG. 21 is a diagram showing a modified embodiment of the configuration illustrated in FIG. 16. FIG. 22 is a diagram showing another modified embodiment of the configuration illustrated in FIG. 16. FIG. 23 is a diagram showing the entire network system in which the information processing apparatus illustrated in FIG. 15 is connected.

#### BEST MODE FOR CARRYING OUT THE INVENTION

The present invention will now be described in detail by way of example with reference to the accompanying drawings.

Referring first FIGS. 1 and 2, the following describes a first preferred embodiment of the present invention. In FIG. 1, there is shown a schematic configuration diagram of an information processing apparatus (101) comprising at least a CPU (102), a main memory (103), and an external storage device (104). A theoretical system bus (114) is extended to connect the CPU (102), a main memory controller (117), and an external storage device controller (115). The main memory controller (117) and the external storage device controller (115) are connected to the main memory (103) and the external storage device (104), respectively. While an actual signal line connection is arranged as shown in FIG. 7, a schematic signal line connection can be represented as shown in FIG. 1 in view of a theoretical data flow.

The CPU (102) comprises a microprocessor (105), a cryptographic processing algorithm ROM (106), a cryptographic processing hardware circuit (107), a RAM (108), a key custody area (112), and an external bus controller (109). These component circuits are integrated on a single semiconductor chip.

In the CPU (102), a microprocessor bus (110) is extended to connect the cryptographic processing algorithm ROM (106), the cryptographic processing hardware circuit (107), the RAM (108), and the external bus controller (109).

09885410-071701

In the present preferred embodiment, data is encrypted within the CPU.

For encryption of a file (111), encryption processing is carried out through the cryptographic processing algorithm ROM (106), using the cryptographic processing hardware circuit (107) if necessary. Key data to be used for encryption may be generated in the CPU (102), or data prepared in advance may be used for encryption. Note that the key data must be retained in the key custody area (112) contained in the CPU (102). In encryption processing, if intermediate result data is generated, the generated intermediate result data is stored in the RAM (108). Thus, an encrypted file (113) is produced from the file (111).

The encrypted file (113) is stored into the external storage device (104) through the system bus (114) and the external storage device controller (115).

For decryption of an encrypted file (116) stored in the external storage device (104), decryption processing is carried out in the reverse order of encryption.

To be more specific, the encrypted file (116) is read into the RAM (108) from the external storage device (104) through the external storage device controller (115). Then, decryption processing is carried out through the cryptographic processing algorithm ROM (106), using the

09885410-07101

cryptographic processing hardware circuit (107) if necessary.

For encrypting/decrypting a large amount data at high speed, common key cryptography (symmetric key cryptography) is employed in which a common key is applied to both encryption and decryption. In the common key cryptography, encryption and decryption are carried out in just the reverse order to each other, i.e., identical minimum-unit processing steps are performed in encryption and decryption. The cryptographic processing algorithm ROM (106) also contains a decryption processing procedure. The cryptographic processing hardware circuit (107) is available for decryption as well as encryption.

In FIG. 2, there is shown a process through which the file (111) indicated in FIG. 1 is produced.

An application program (201) is retained in the external storage device except when it is activated. When the application program is initiated, the application program is loaded into the main memory and put in an operable state (202). Then, the application program in the operable state (202) issues a request for working area allocation to an operating system serving as a control program for the information processing apparatus. In answer to this request, the operating system for controlling the

information processing apparatus allocates a memory space of the RAM (108) as a working area (203).

In the above-mentioned state, the application program (202) is executed under control of the microprocessor (105), and information is created and stored into the working area (203). Of the information thus created and stored in the working area (203), data to be stored into the external storage device is produced as a file (111).

The application program (202) itself resides in the main memory. In allocation of the working area (203) on the RAM (108) for execution of the application program, a memory control function of the microprocessor is used under the aegis of the operating system for controlling the information processing apparatus. Through the use of the memory control function of the microprocessor, logical addresses representing the working area for the application program are converted into physical addresses in the RAM (108).

The key custody area (112) may be provided in the RAM (108), but it is required that the key custody area should be of a nonvolatile memory type. Instead of the RAM (108), a nonvolatile ROM such as an EEPROM or flash ROM may be used to provide the key custody area. It is also preferable to use a battery-backed SRAM in lieu of the RAM (108). Where the battery-backed SRAM is used to provide the key custody

09889410-071701



for controlling the information processing apparatus allocates a memory space of the RAM (108) as a working area (303). In this state, the application program (302) is executed under control of the microprocessor (105), and information is created and stored into the working area (303). Of the information thus created and stored in the working area (303), data to be stored into the external storage device is produced as a file (111).

The file (111) thus produced is encrypted through the cryptographic processing algorithm ROM (106), using the cryptographic processing hardware circuit (107) if necessary. Thus, an encrypted file (113) is produced from the file (111). The encrypted file (112) is then stored as an encrypted file (116) into the external storage device.

Although main memory disposed outside the CPU is not shown in FIG. 3, there may also be provided such a modified arrangement that application programs used for generating sensitive information and any other application programs are handled differently. To be more specific, the application programs used for generating sensitive information are carried out in the RAM (108), and the other application programs are carried out on the main memory disposed outside the CPU in a conventional manner.

Where the RAM (108) is used as the main memory as described above, a transfer of non-encrypted sensitive



information takes place outside the CPU (102) only when the application program (301) is loaded into the RAM (108), and the application program itself is executed in a secure manner.

With reference to FIG. 4, the following describes a third preferred embodiment of the present invention.

In the third preferred embodiment, an encrypted application program (401) is retained in the external storage device (104). The encrypted application program is decrypted in the CPU of the information processing apparatus. Hence, a decrypted form of the application program is not output onto the bus (114), i.e., the decrypted application program is handled within the CPU in a closed fashion. This prevents any other person from tapping the application program.

The third preferred embodiment is explained in detail below. At the time of startup of the encrypted application program (401) retained in the external storage device, the encrypted application program (401) is transferred to the RAM (108) in the information processing apparatus over the bus (114). Through this transfer, the RAM (108) is loaded with the encrypted application program (402). Then, in the RAM (108), the encrypted application program (402) is decrypted in the RAM (108) to take an original form of an application program (403). In this state, the application

09859410-071701

program (403) is run to generate information using a working area (404) in the RAM (108). Then, any necessary parts are selected from the generated information to produce a file (111). The file (111) thus produced is encrypted to produce an encrypted file (113). The encrypted file (113) is stored as an encrypted file (116) into the external storage device (104).

As explained above, an application program itself is stored as an encrypted file into the external storage device, thereby enhancing security protection further.

In the third preferred embodiment, for generating the encrypted application program (401), the application program itself is encrypted as a file (111).

Then, with reference to FIGS. 5 and 6, the following describes the external bus controller in the present invention.

The external bus controller (109) used in each of the first to third preferred embodiments is designed to control data input/output between the inside of the CPU and the outside thereof. For example, in access from the microprocessor (105) to the cryptographic processing algorithm ROM (106), cryptographic processing hardware circuit (107) or RAM (108) for execution of cryptographic processing, the external bus controller performs control so that an access signal will not go out of the CPU. In a case

09889410-071701

where security protection is not affected when an access signal from the microprocessor (105) is output from the CPU, the external bus controller may allow output of the access signal from the CPU. Information allowed to be output from the CPU includes data which is to be transferred to another information processing apparatus without being subjected to encryption processing.

As shown in FIGS. 5 and 6, an external bus controller (501) corresponding the external bus controller (109) is disposed between a set of three internal buses from a microprocessor (502) and a set of three external buses extending out of the CPU. The three internal buses are a control bus (503), an address bus (504), and a data bus (505), and the three external buses are an external control bus (506), an external address bus (507), and an external data bus (508). The external bus controller (501) comprises an external control bus generator (509), an address comparator (510), an address direction controller (512), a data direction controller (513), a mask signal generator (511), and signal mask circuits (514) and (519).

A bus cycle start signal, a bus direction indicating signal, a bus cycle end signal, a bus arbitration signal, and other control signals from the microprocessor are fed through the control bus (503) and the external control bus (506). Bus cycle control is carried out by these signals.

09889410-071701

09889410-071701

The external control bus generator (509) monitors the bus cycle start signal, bus direction indicating signal, bus cycle end signal, bus arbitration signal, and other control signals from the microprocessor. In the external control bus generator (509), it is determined whether the microprocessor has a bus access right or not. Information on the result of judgment is given to the address direction controller (512), and the same information is also given to the address comparator (510). The address comparator (510) has information regarding addresses allocated to the cryptographic processing algorithm ROM (106), cryptographic processing hardware circuit (107), and RAM (108) in the CPU (102), and in the address comparator (510), these addresses are compared with an address input through the address bus (504) or the external address bus (507).

If the external control bus generator (509) judges through the control bus (503) that the microprocessor has the bus access right, the address comparator (510) checks an address signal from the microprocessor. Then, if it recognized that access to an address of the RAM (108) is attempted, the address comparator (510) notifies the external control bus generator (509) thereof so that an external bus control signal will not be driven. The address comparator (510) also notifies the mask signal generator (511) thereof so as to output a mask signal to the signal

mask circuits (514) and (519). Thus, the external address bus (507) and the external data bus (508) are controlled to remain inactive. Alternatively, an address value or a data value is forcedly made invariable.

If the external control bus generator (509) judges through the control bus (503) that the microprocessor does not have the bus access right, the address comparator (510) checks the external address bus. Then, if it is detected that access to an address of the RAM (108) is attempted, the address comparator (510) notifies the external control bus generator (509) thereof. The external control bus generator (509) conducts control not to indicate a bus cycle to the control bus (503). Alternatively, a mask signal is output to the signal mask circuits (514) and (519) so that the address bus (504) and the data bus (505) will not be driven, or an address value or a data value is forcedly made invariable.

In implementation of making an address value or a data value invariable, gate logic is changed as in a gate (602) of a signal mask circuit (601) and a gate (604) of a signal mask circuit (603) shown in FIG. 6.

Thus, through the use of the address signal mask circuit, it is possible to make address conversion to indicate any area other than the RAM (108) area, i.e., an

area where no problem will arise in read and write operations.

As described above, analysis of internal processing in the CPU (102) by tapping the system bus (114) connected to the CPU (102) can be made infeasible. This enhances the security of cryptographic processing to be carried out in the CPU (102).

Then, with reference to FIGS. 7 to 13, the following describes how key information is handled.

Key information is required for encryption and decryption, and the secrecy of the key information is critical to security protection of the entire system. In conventional practice, key information is furnished from an external source, and the key information furnished externally must be strictly controlled by an individual person for ensuring system security.

According to the present invention, key information necessary for encryption is generated in a semiconductor device, and the key information thus generated is retained in the semiconductor device only. When the key information needs to be output from the semiconductor device, it is delivered using means recognizable only by a particular party. Random numbers are used for generation of key information. In logic-based generation of random numbers, pseudo-random numbers are generated commonly. That is to

09889410-071704

FIG. 7 shows an exemplary configuration of the random number generator (118) for generating random numbers through the use of a physical phenomenon. The random number generator (118) shown in FIG. 7 generates random numbers using noise produced by a constant-voltage diode or a Zener diode. As illustrated in FIG. 7, the random number generator (118) comprises a low-pass filter (704) including a constant-voltage diode (701), a resistor (702) and a capacitor (703), a comparator (705), and a flip-flop (706).

The constant-voltage diode (701) produces noise having a signal waveform (707). This noise is caused by the physical phenomenon that an avalanche breakdown occurs randomly at a semiconductor junction inside the constant-voltage diode (701). The noise thus produced is

fed through the low-pass filter (704) to yield a signal waveform (708) which has a value approximating to the average value of the signal waveform (707). These two signal waveforms (708) and (707) are input to the comparator (705) for conversion into a binary signal having random pulse widths such as a signal waveform (709). The binary signal thus attained is further applied to the flip-flop (706), in which it is synchronized with a reference clock signal available in the semiconductor device. Thus, a random-bit signal waveform (710) is provided.

Random numbers are generated by inputting a necessary length of the random-bit string to a shift register or by counting the number of random pulses in a time unit.

Non-reproducible random numbers can thus be attained without using seed information which is required in common practice of random number generation. Further, through the low-pass filter (704), the average value (708) of the signal containing noise (707) is determined, and the average value thus determined is compared with the signal containing noise. Thus, in the random number generator, an adverse effect on random number generation can be prevented even if voltage fluctuations occur due to such a cause as temperature variation in the constant-voltage diode.

While the constant-voltage diode is used as a noise source in the preferred embodiment illustrated in FIG. 7,

05885410-071701



it is to be understood that the present invention is not limited to the use of the constant-voltage diode and that any circuit producing noise based on a physical phenomenon may be used instead thereof.

Referring to FIG. 8, there is shown an exemplary configuration of the key custody area for storing generated key information in secrecy. The key custody area (112) shown in FIG. 8 is an example of a battery-backed SRAM configuration.

The CPU (102) of the present invention is divided into the following sections; an SRAM (804), an SRAM control circuit (809), and a CPU internal logic block (802) for other circuits. A power supply (805) dedicated for the SRAM (804) and the SRAM control circuit (809), and a main power supply (803) dedicated for the internal logic block (802) are provided. The main power supply (803) and the power supply (805) for the key custody area (112) are connected to the SRAM (804) through diodes (806) and (807) to furnish a power (808) thereto. The power (808) is also furnished to the SRAM control circuit (809). A gate (810) monitors a reset signal (811) and the main power supply (803) used for the internal logic block, and the gate (810) fixedly nullifies all the signals to the SRAM (804) until initialization of the internal logic block is completed with the main power supply (803) turned on. Thus, even under condition that only the

key custody area (112) is supplied with power while the other circuits are not supplied with power, it is possible to eliminate an excess leakage current. Further, even if a noise is applied to any part of the circuits not supplied with power or even if a malfunction occurs due to a voltage lower than a guaranteed level, an adverse effect can be shut off. When the gate (810) outputs a "Low" signal, a gate (812) outputs a "Low" signal to provide no change regardless of a voltage level on address signal lines (813). Further, since a gate (815) outputs a "Low" signal, an output impedance of a buffer (816) increases to prevent current leakage to data signal lines (818). Still further, since a gate (820) outputs a "Low" signal, an output impedance of a buffer (821) increases to prevent any data from going to a data signal line (819) regardless of a voltage on data signal lines (823). Still further, a gate (824) outputs a "High" signal to nullify a control signal (826), thereby stopping operation of the SRAM. Further, since gates (810), (812), (815), (820) and (824) and a buffer (821) are formed using CMOS-structure elements, leakage currents to input signals (813), (817), (823), (822) and (825) can be suppressed to a minimum. Therefore, even if power to the internal logic block (802) is suspended, there occurs no leakage of power from the key custody area power supply (805). Thus, power consumption of the key custody area power

05889410-071701

supply (805) can be restricted to a required minimum, thereby making it possible to prolong the service life of a backup battery (828).

When power is furnished from the main power supply (803) and initialization of the internal logic block (802) is completed, the gate (810) outputs a "High" signal. Then, the gate (812) allows data on the address signal line (813) to flow to the address signal lines (814). Under condition that an SRAM read signal (817) is effective, the gate (815) enables the buffer (816) for allowing data on the data signal line (819) to go to a data signal lines (818). Under condition that an SRAM write signal (822) is effective, the gate (820) enables the buffer (821) for allowing data on the data signal lines (823) to flow to the data signal lines (819). The gate (824) allows data on the control signal line (825) to flow to the control signal line (826). Through the above-mentioned sequence, the SRAM (804) can be accessed normally.

Moreover, a variety of sensors are provided on a case or housing in which the CPU (102) is accommodated, and an alarm detector (827) receiving signals from these sensors is provided for controlling the key custody area power supply (805) connected with the battery (828). Upon detection of an abnormality such as disassembling/removing the case or housing, power to the SRAM (804) is stopped to

09889410-07101

delete key information. Further, under condition that the main power supply (803) is turned on, even if the key custody area power supply (805) is shut off by actuation of the alarm detector (827), power is furnished from the main power supply (803) to the SRAM (804). Therefore, an alarm detection signal (829) is input to the internal logic block (802) to inform occurrence of an abnormality, thus restricting or stopping operation of the CPU (102).

In a modified embodiment, the power supplies to the SRAM (804) may be integrated in the CPU (102). Further, there may also be provided an arrangement wherein the main power supply and battery power supply are integrated in the alarm detector (827) and wherein the power supplies to the SRAM (804) are shut off upon detection of any abnormality.

In terms of circuit mounting, it is required to provide each of the signal and power lines between the alarm detector (827) and the CPU (102) along the shortest possible route in an inner layer of a substrate therefor so that probing cannot be made with ease. Further, it is required to provide such a security protective means as connecting a plurality of signal lines along different wiring paths. Thus, even if the information processing apparatus is disassembled, cipher key information cannot be extracted out of the semiconductor device.

In the present invention, two kinds of cipher keys are generated for different purposes; a key necessary for identifying the semiconductor device concerned, and a key used for encrypting information. The former is an authentication key, and the latter is an information enciphering key. A key to be generated frequently is an information enciphering key, i.e., in principle, an information enciphering key is generated each time information is encrypted. An authentication key is generated at a predetermined interval of time, e.g., per month or per year, for use as information for identifying the semiconductor device concerned.

Referring to FIG. 9, there is shown how enciphering and deciphering keys are handled in a situation where the same information processing apparatus is used for encryption and decryption.

Information (901) produced in the CPU (102) semiconductor chip is encrypted and stored as an encrypted file (116) into the external storage device (104) or the like. Thereafter, in decryption for using the contents of the encrypted file in the CPU (102) again, it is just required that key information (902) should reside in the CPU (102) only. For making it possible to handle the encrypted files (113) and (116) in the CPU (102) only, the key information (902) is generated by the random number

Further, in a case where different cipher keys (904) and (906) are used for encrypting a plurality of information (903) and (905) respectively, it is required to store key information (904) and (906) into the key custody area (112).

Referring to FIG. 10, the key custody area (112) holds just a key (1001) generated by the random number generator (118). In encryption of information (1002) or (1003), a key (1004) or (1005) generated in association therewith is encrypted with the key (1001) to produce a cipher key (1006) or (1007). The information (1002) or (1003) is encrypted with the cipher key (1004) or (1005) thus produced to create an encrypted file (1008) or (1009). The encrypted file (1008) thus created and the cipher key (1006) are collectively stored as a file (1010) into the external storage device, and the encrypted file (1009) thus created and the cipher key (1007) are collectively stored as a file (1011) into the external storage device. In this fashion,

the amount of key information to be held in the key custody area (112) can be reduced.

Referring to FIG. 11, there is shown how enciphering and deciphering keys are handled in a situation where different apparatuses are used for encryption and decryption respectively. In this case, it is required to ensure that each communicating party is a correct one. That is to say, it is required to authenticate the identity of each communicating party.

As a means for authenticating the identity of the communicating party, asymmetric key cryptography is employed. In the asymmetric key cryptography, different keys are used for encrypting information to provide a cipher text and for decrypting the cipher text to attain the original information. The asymmetric key cryptography is also referred to as public key cryptography, in which enciphering and deciphering keys are used. One of the enciphering and deciphering keys is made public, and the other is kept secret. Information encrypted with the enciphering key can be decrypted only with the deciphering key corresponding thereto. One of the two keys which is made public is referred to as a public key, and the other one which is kept secret is referred to as a private key. Cipher text information encrypted with the public key can be decrypted only with the private key corresponding thereto, and cipher

text information encrypted with the private key can be decrypted only with the public key corresponding thereto. In this cryptographic technique, cipher text information can be transmitted in a form decipherable only by a particular party, and the identity of a sender can be authenticated.

For sending information only to a particular recipient, a sender attains a public key from the recipient first. Then, using the public key of the recipient, the sender encrypts information to be sent. The information thus encrypted is a cipher text which can be decrypted only with a private key kept secret at the recipient, not with the same public key. In this manner, the encrypted information can be sent to the particular recipient securely. In general, since the asymmetric key cryptography (public key cryptography) requires a substantial amount of time due to complex processing thereof, the following cryptographic processing is implemented in most actual applications: Information to be sent is encrypted by symmetric key cryptography (common key cryptography) in which cipher key information is generated through random number generation each time. Using asymmetric key cryptography (public key cryptography), the cipher key information thus generated is sent to a recipient in secrecy.

09859410.07101



For authenticating the identity of a sender, the sender encrypts information itself or any message corresponding thereto (e.g., message digest) with a private key and then sends the encrypted information to a recipient. The recipient attains a public key of the sender, and decrypts the encrypted information with the public key of the sender. In authenticity check, if it is judged that the information has been encrypted with the private key owned only by the sender, the recipient can authenticate the identity of the sender.

In an apparatus A (1101a) and an apparatus B (1101b) shown in FIG. 11, public keys (1104a) and (1104b) and private keys (1105a) and (1105b) are generated respectively as apparatus identification information in advance. In a public key cryptographic technique based on modulo arithmetic, each of these keys is generated using the product of two prime numbers. For this purpose, random number generation (1102a) (1102b) is performed to generate random numbers. Then, it is checked whether the random numbers thus generated are prime numbers or not. Random prime numbers can thus be attained (1103a) (1103b). The generated keys are used for identifying the semiconductor devices concerned.

The following describes a situation in which information (1116) is sent from the apparatus A (1101a) to the apparatus B (1101b).

In transfer of the information (1116) from the apparatus A (1101a) to the apparatus B (1101b), the apparatus A (1101a) encrypts the information (1116) with a key so that the information (1116) can be decrypted only by the apparatus B (1101b). It is desirable to use a key effective only for transfer of the information (1116) and another key for transfer of different information. In this usage of keys, even if key information leaks, possible damage can be kept to a minimum. For this purpose, a common key (1111) for encrypting information (1116) to be generated each time must be sent only to the apparatus B.

In implementation of the above-mentioned arrangement, the apparatus A (1101a) issues a request for public key transfer (1106). In response to this request, the apparatus B (1101b) transfers the public key (1104b) to the apparatus A (1101a). Then, in the apparatus A (1101a), random numbers are generated (1109), and a common key (1111) is generated using the random numbers thus generated (1110). Using the generated common key (1111), a public key of the apparatus B (1108) received from the apparatus B (1101b) is subjected to public key encryption (1112) to produce encrypted key information (1113). Further, using the

09889440-074701

generated common key (1111), the information (1116) is subjected to common key encryption (1115) to produce encrypted information (1117). The encrypted key information (1113) and encrypted information (1117) thus produced are sent to the apparatus B. In this manner, the information (1116) can be sent to the apparatus B in a form decipherable only by the apparatus B. In the apparatus B (1101b), the encrypted key information thus received (1119) is subjected to public key decryption (1120) using the private key (1105b) of the apparatus B to attain a common key (1121). Further, using the common key (1121), the encrypted information received as mentioned above (1122) is subjected to common key decryption (1123) to attain information (1124).

Moreover, for proving that information transfer (1118) has been made from the apparatus A (1101a), a message digest of the information (1116) is determined as a hash value (1202) using a hash function (1201) as shown in FIG. 12. Then, the hash value (1202) thus determined is subjected to public key encryption (1203) using the private key (1105a) of the apparatus A to generate an encrypted hash value (1204). The public key (1104a) of the apparatus A is transferred (1205) to the apparatus B (1101b), and the encrypted hash value (1204) is also transferred to the apparatus B (1101b) as a signature of the apparatus A

09889410-071701

(1101a). In the apparatus B (1101b), using the public key of the apparatus A received from the apparatus A (1207), the encrypted hash value received therefrom (1208) is subjected to public key decryption (1209) to attain a hash value (1210) corresponding to the hash value generated by the apparatus A (1101a). On the other hand, using a hash function (1211), a hash value (1212) is determined from the information (1124) received from the apparatus A (1101a). The two hash values (1210) and (1212) mentioned above are compared (1213) to check whether they are identical or not. If a match is found therebetween, it can be verified that the sender of the information (1124) is the apparatus A (1101a).

While a hash value of the information (1116) is determined in the arrangement exemplified in FIG. 12, there may also be provided such a modified arrangement that the information itself (1116) is encrypted with the private key (1105a) of the apparatus A and the encrypted information is transferred to the apparatus B (1110b) together with the public key (1104a) of the apparatus A. In particular, this modified arrangement is preferable in a case where the size of the information (1116) is relatively small.

In acquisition of a public key of a desired communicating party, the public key may be attained directly from the communicating party as illustrated in the above-mentioned preferred embodiments, or the public key of

the communicating party may be attained through a third party having no interest in relation with the communicating party.

As to a public key attained from each communicating party, it is necessary to ensure that the attained public key is authentic, i.e., the communicating party is not an imposter.

Referring to FIG. 13, there is shown an arrangement for checking whether the public key received from each communicating party in the preferred embodiments shown in FIGS. 11 and 12 is authentic or not. In the arrangement shown in FIG. 13, an apparatus C (1301) is provided as a certification authority server for authenticating the identity of each terminal apparatus. The apparatus C (1301) authenticates a public key of each apparatus subscribing to the communication system concerned. For this purpose, in the apparatus C (1301), random numbers are generated (1302), and from the generated random numbers, prime numbers are taken (1303). Then, using the prime numbers, a public key (1304) and a private key (1305) of the apparatus C are generated. The private key of the apparatus C must be kept in absolute secrecy for ensuring the security of the communication system concerned.

In the apparatus A (1101a) and the apparatus B (1101b), the public keys (1104a) and (1104b) and the private

keys (1105a) and (1105b) are generated for apparatus identification as mentioned before. Then, each of the public keys is transferred to the apparatus C (1301) to make a request for certification (1316a) (1316b). Upon receiving the request for certification, the apparatus C (1301) processes the public key (1306a) (1306b) of each of the apparatuses A and B through public key encryption (1307a) (1307b) using the private key (1305) of the apparatus C to produce a certificate (1308a) (1308b). The certificate thus produced and the public key (1304) of the apparatus C are combined to provide the result of certification (1309a), which is then transferred to each of the apparatuses A and B (1317a) (1317b).

In each of the apparatuses A and B, the result of certification, i.e., a certificate of the public key thereof is stored. When the apparatus B receives a request for public key transfer from the apparatus A for sending information, the apparatus B transfers the public key (1105b) thereof and the certificate thereof to the apparatus A. The certificate thus transferred indicates that the apparatus B has been certified by the apparatus C. When the apparatus A receives the certificate (1312), public key decryption (1313) is performed using the public key of the apparatus C stored in the apparatus A. Then, the public key (1314) of the apparatus B is extracted from the certificate

(1312), and the public key thus extracted is compared with the public key (1108) transferred from the apparatus B (1101b). In this manner, the authenticity of the public key of the apparatus B can be verified.

In certification by the apparatus C using the public key of each of the apparatuses A and B, electronic verification and meticulous inspection are carried out to make sure that each of the apparatuses A and B is not modified or tampered by any third party.

Through the above-mentioned procedure, the authenticity of the public key of each communicating party can be ensured.

Then, with reference to FIG. 14, the following describes a fourth preferred embodiment of the present invention.

In FIG. 14, there is shown a schematic configuration of a typical information processing apparatus. An information processing apparatus (1401) comprises a plurality of semiconductor devices. A CPU (1402) is connected to a cache memory and a main memory controller (1405) through a processor bus (1404). The main memory controller (1405), including a system bus controller, is provided with a memory bus (1413) and a system bus (1407). The memory bus (1413) is extended to a main memory (1406), and the system bus (1407) is extended to an external storage

device controller (1408), a display circuit controller (1410), a communication circuit controller (1411), and a peripheral I/O controller (1412). The display circuit controller (1410) may be connected to the main memory controller including the system bus controller (1405). The external storage device controller (1408) is connected to an external storage device (1409).

Since an address region of the main memory (1406) is different from that of each component connected to the system bus (1407), the main memory controller including the system bus controller (1405) conducts a changeover in address region access.

In the information processing apparatus (1401) which can be regarded as a system, the CPU (1402) serves as a main processor presiding over operations in the system. According to the present invention, cryptographic processing is carried out in the CPU in a closed fashion. For example, as shown in FIG. 1, the CPU (1402) comprises the microprocessor (105), cryptographic processing algorithm ROM (106), cryptographic processing hardware circuit (107), RAM (108), key custody area (112), and external bus controller (109), which are all integrated on a single semiconductor chip. Further, according to the present invention, an information processing apparatus

09889440-071701



comprising a plurality of CPUs may be used as shown in FIGS. 19 and 20.

With reference to FIG. 15, the following describes a fifth preferred embodiment of the present invention.

In the fifth preferred embodiment shown in FIG. 15, an information processing apparatus (1501) is connected to another information processing apparatus for communication therewith. In lieu of the external storage device indicated in FIG. 1, a communication circuit controller is provided in the fifth preferred embodiment. The communication circuit controller may be disposed outside the information processing apparatus (1501).

The information processing apparatus (1501) comprises a CPU (1502) and a communication circuit controller (1503), which are interconnected through a system bus (1514). The CPU (1502) includes a microprocessor (1505), a cryptographic processing algorithm ROM (1506), a cryptographic processing hardware circuit (1507), a RAM (1508), an external bus controller (1509), and a key custody area (1512). These component circuits are interconnected through a microprocessor bus (1510).

While the information processing apparatus shown in FIG. 15 comprises the CPU and the communication circuit controller, other components such as a main memory and an external storage device may be included in the information



for high-speed encryption/decryption of a large amount of data. In the common key cryptography, since the same key is used for encryption and decryption, the information processing apparatuses A and B must have an identical key. The identical key may be set in advance in each of the information processing apparatuses A and B. Alternatively, before transmission of encrypted data, mutual authentication may be made between the information processing apparatuses A and B in a fashion that an enciphering key is shared therebetween. For the mutual authentication, cryptographic processing is carried out in the CPU.

FIG. 23 shows a scheme in which the information processing apparatuses A and B are connected over a network.

In the RAM (1508), encrypted data is re-edited per communication, and according to a communication protocol, the re-edited data is transferred to the communication circuit controller (1503) to enable secure communication. There may be provided such a modified arrangement as mentioned below: Data encrypted in the RAM (1508) is transferred to the communication circuit controller (1503), and in the communication circuit controller (1503), the encrypted data is re-edited per communication. Then, according to a communication protocol, the re-edited data is sent over the communication line (1504).

With reference to FIGS. 16, 17, 18, 21 and 22, the following describes a sixth preferred embodiment of the present invention.

In FIG. 16, there is shown a configuration in which a group of external storages including a magnetic disk (1601) is controlled by a disk system controller (1602). The disk system controller (1602) is connected to a host system (1603) which is a higher-ranking information apparatus.

The magnetic disk (1601) contains data stored as a file and file location information indicating where the file is stored on the magnetic disk. In a small-type information processing apparatus such as a personal computer, a file system program for managing files and file location information may be carried out by a CPU contained in the small-type information processing apparatus. In a disk system controller designed for implementing high-speed operations with high reliability, the disk system controller itself may manage files and file location information.

The sixth preferred embodiment of the present invention is particularly suitable for the latter case. In the host system, a file (1604) and a file identifier (1605) are managed. Whether the file (1604) is encrypted or not depends on the host system, i.e., it is not required for the disk system controller (1602) to be concerned therewith. In

09889410-07101

the disk system controller (1602), file location information (1606) on the magnetic disk (1601) is encrypted for management.

In the sixth preferred embodiment of the present invention, operations to be performed until the host system reads out an encrypted file (1607) are as follows:

First, the host system sends a file identifier (1605) corresponding to a necessary encrypted file to the disk system controller (1602), thereby indicating a request for reading out the encrypted file. Upon receipt of this request, the disk system controller (1602) reads the encrypted file location information (1606) out of the magnetic disk (1601). In the disk system controller (1602), the encrypted file location information (1606) is decrypted to extract file location information (1608). Then, the disk system controller (1602) searches the extracted file location information (1608) for the file identifier (1605) to attain actual file location information. Using the file location information thus attained, the disk system controller (1602) reads the necessary encrypted file (1607) out of the magnetic disk (1601), which is then transferred to the host system (1603).

With reference to FIG. 17, the following describes how a file is written onto the magnetic disk. Until the file location information (1608) is extracted, the same

operations as those in the case of encrypted file reading mentioned above are carried out. Using the extracted file location information (1608), a free space of the magnetic disk (1601) is checked, and then the encrypted file (1604) is written into the free space of the magnetic disk (1601). After completion of writing the encrypted file (1604), the file location information (1608) is updated and encrypted. Thereafter, the encrypted file location information (1701) is written onto the magnetic disk (1601).

With reference to FIG. 18, the following describes a configuration of a disk system controller (1801) according to the present invention.

The disk system controller (1801) of the present invention includes a CPU (1802) serving as an internal disk system processor, a magnetic disk interface (1813), and a host system interface (1814). The CPU (1802) comprises a microprocessor (1805), a cryptographic processing algorithm ROM (1806), a cryptographic processing hardware circuit (1807), a RAM (1808), a key custody area (1811), an external bus controller (1809), and a random number generator (1820).

As shown in FIGS. 21 and 22, a plurality of magnetic disk devices may also be connected to an information processing apparatus.

Using a disk system controller such as mentioned above, all the information stored in a magnetic disk can be encrypted to ensure security protection in information storage.

The cryptographic processing hardware circuit of the present invention is a dedicated hardware component for common key cryptography in which a common key is applied to both encryption and decryption. The cryptographic processing hardware circuit comprises a rotator, an adder, an arithmetic logic unit, and other elements. It is also practicable to use such a common key cryptographic method as Multi-series cryptography or M6 cryptography in which bit rotation, addition, and logic operation are primarily performed in units of a certain data length.

In a case where public key cryptography is used, there is provided a modulo arithmetic circuit capable of handling higher computational complexity as a dedicated hardware circuit.

#### INDUSTRIAL APPLICABILITY

According to the present invention, cryptographic processing can be carried out without outputting sensitive information onto a system bus and processor bus in an information processing apparatus. Since cryptographic processing, sensitive information to be





CLAIMS

1. An information processing apparatus including,  
a control device for performing predetermined processing of information, and  
a bus for interconnecting said control device and other component devices of said information processing apparatus,  
characterized in that said control device is integrated on a single semiconductor chip in which key information is generated and in which sensitive information is encrypted.
2. An information processing apparatus as claimed in claim 1,  
wherein said control device comprises an external bus controller for preventing non-encrypted sensitive information from being output onto said bus.
3. An information processing apparatus as claimed in claim 2;  
wherein information not requiring encryption is output onto said bus through said external bus controller.

09885410-071701

4. An information processing apparatus as claimed in claim 1,

wherein a memory device is provided for storing information encrypted by said control device.

5. An information processing apparatus as claimed in claim 1,

wherein said control device comprises means for decrypting encrypted information at an information write operation.

6. An information processing apparatus as claimed in claim 5,

wherein said information processing apparatus is connected to a different information processing apparatus through a network, and

wherein said information processing apparatus decrypts encrypted information which is received from said different information processing apparatus.

7. An information processing apparatus as claimed in claim 1,

wherein a plurality of said processing devices are provided, and cryptographic processing is carried out in each of said processing devices.

09889410-07101

8. An information processing apparatus as claimed in claim 1,

wherein said processing device comprises means for receiving an encrypted program and for carrying out decryption thereof.

9. An information processing apparatus as claimed in claim 1

wherein said processing device comprises:

a microprocessor for carrying out said predetermined processing;

a cryptographic algorithm memory device for storing an algorithm for information cryptographic processing;

a cryptographic processing device for carrying out cryptographic processing according to said algorithm; and

a microprocessor bus for interconnecting said microprocessor, said cryptographic processing algorithm memory device, and said cryptographic processing device.

10. A disk system controller, including a processing device for carrying out information processing to control a magnetic disk for storing encrypted information,

characterized in that, upon receipt of a request for reading out said encrypted information, key information is

0385410-071701

generated, encrypted file location information indicating a location of information stored on said magnetic disk is read out of said magnetic disk, said encrypted file location information thus read out is decrypted in a semiconductor chip on which said processing device is integrated, and according to the decrypted file location information, said encrypted information is read out of said magnetic disk.

11. A disk system controller as claimed in claim 10,  
wherein said disk controller is connected to a plurality of magnetic disks.

12. A disk system controller as claimed in claim 10,  
wherein said disk system controller is connected to an information processing apparatus, and  
wherein said disk system controller reads out said encrypted information from said magnetic disk upon receipt of a request therefrom.

\* \* \* \* \*

05889410-071701

FIG.1

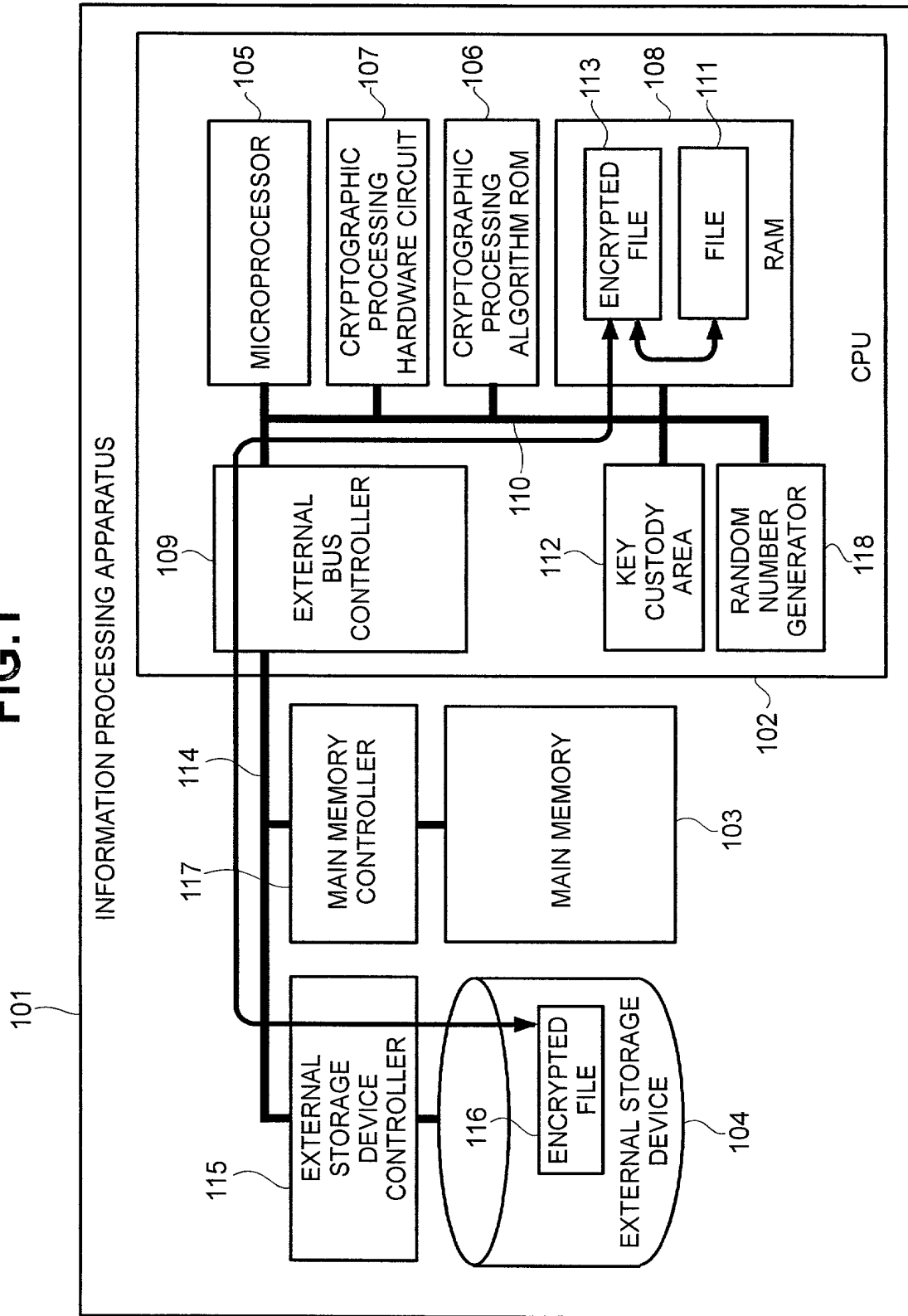


FIG.2

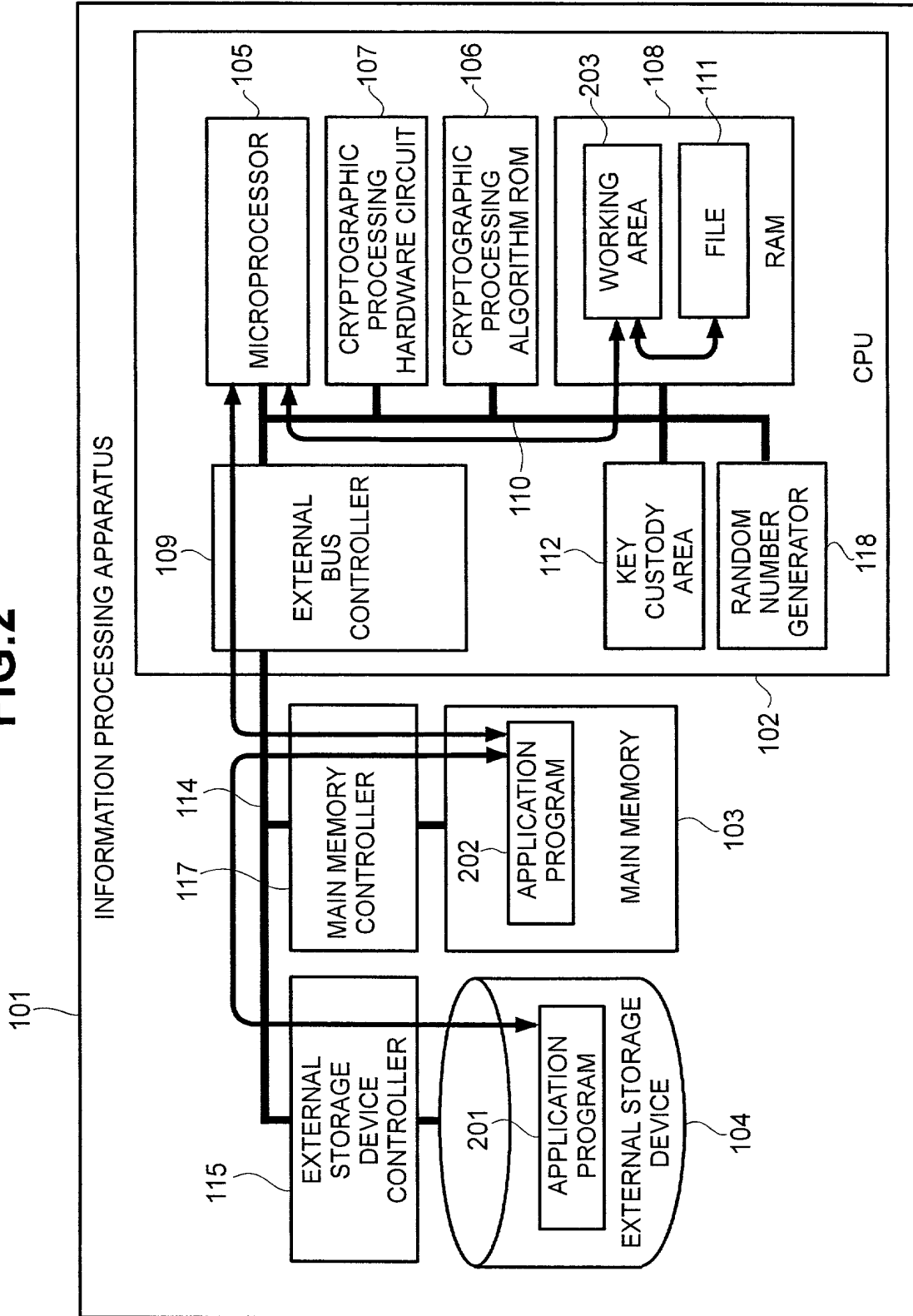


FIG. 3

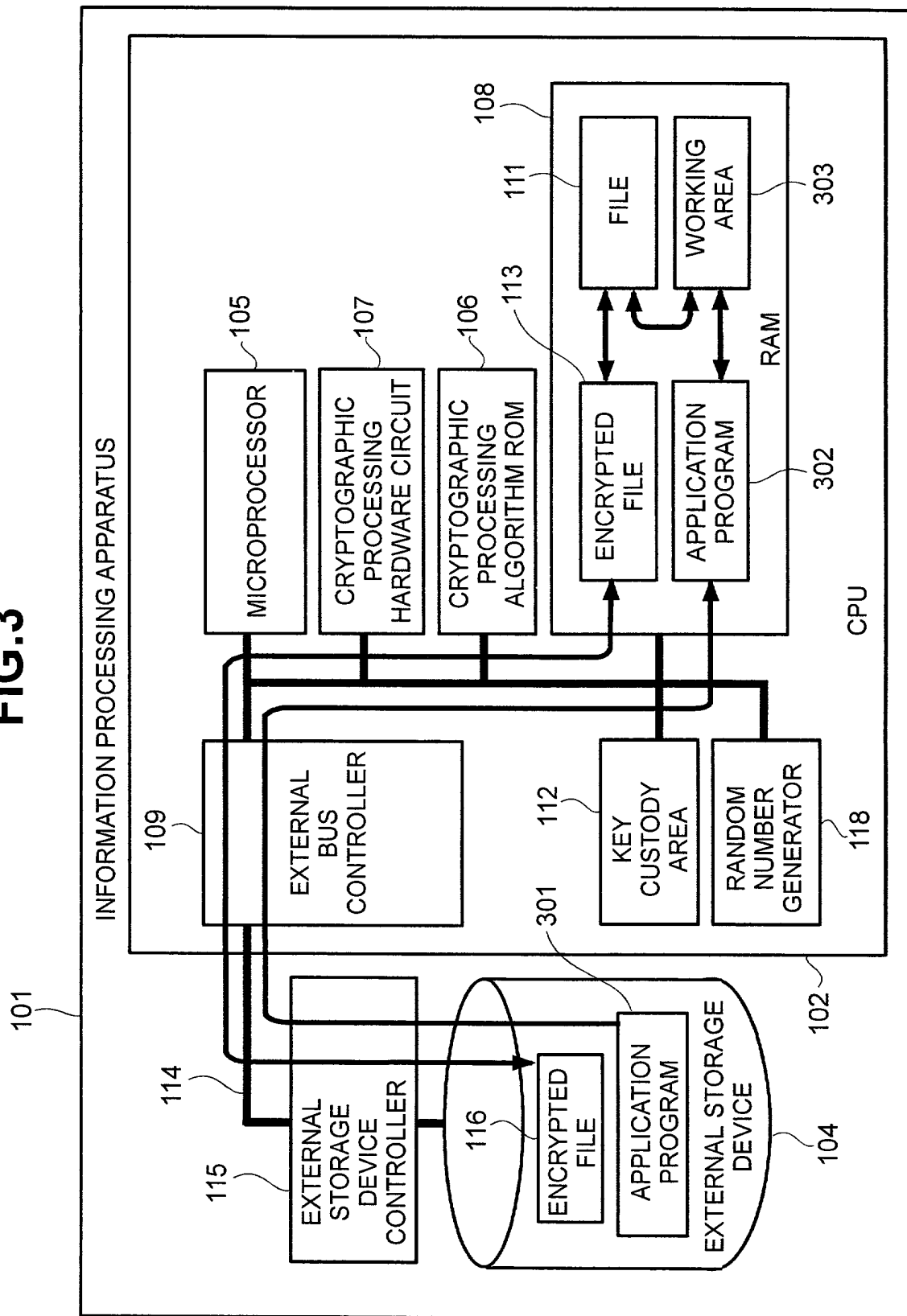


FIG. 4

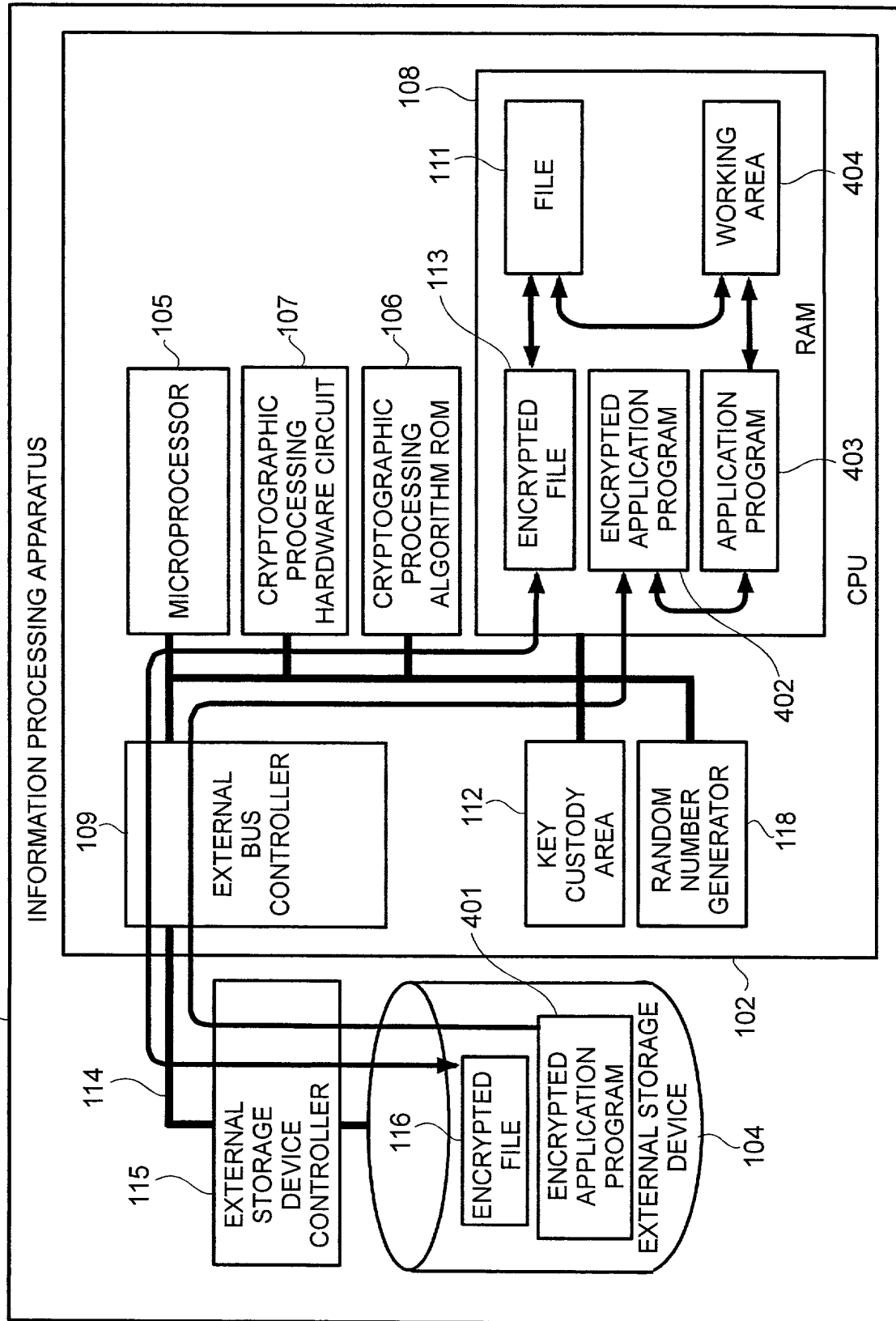
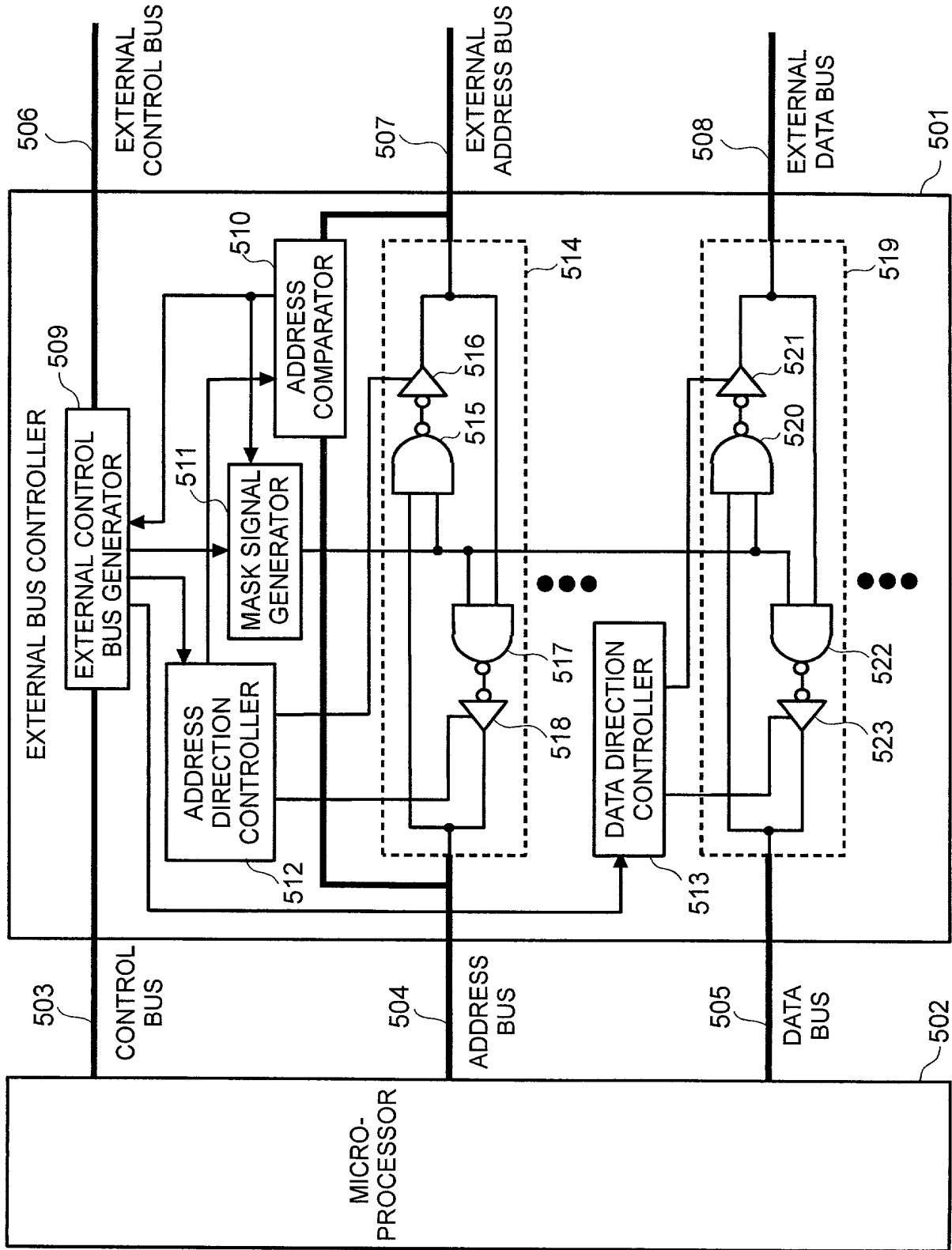




FIG. 5



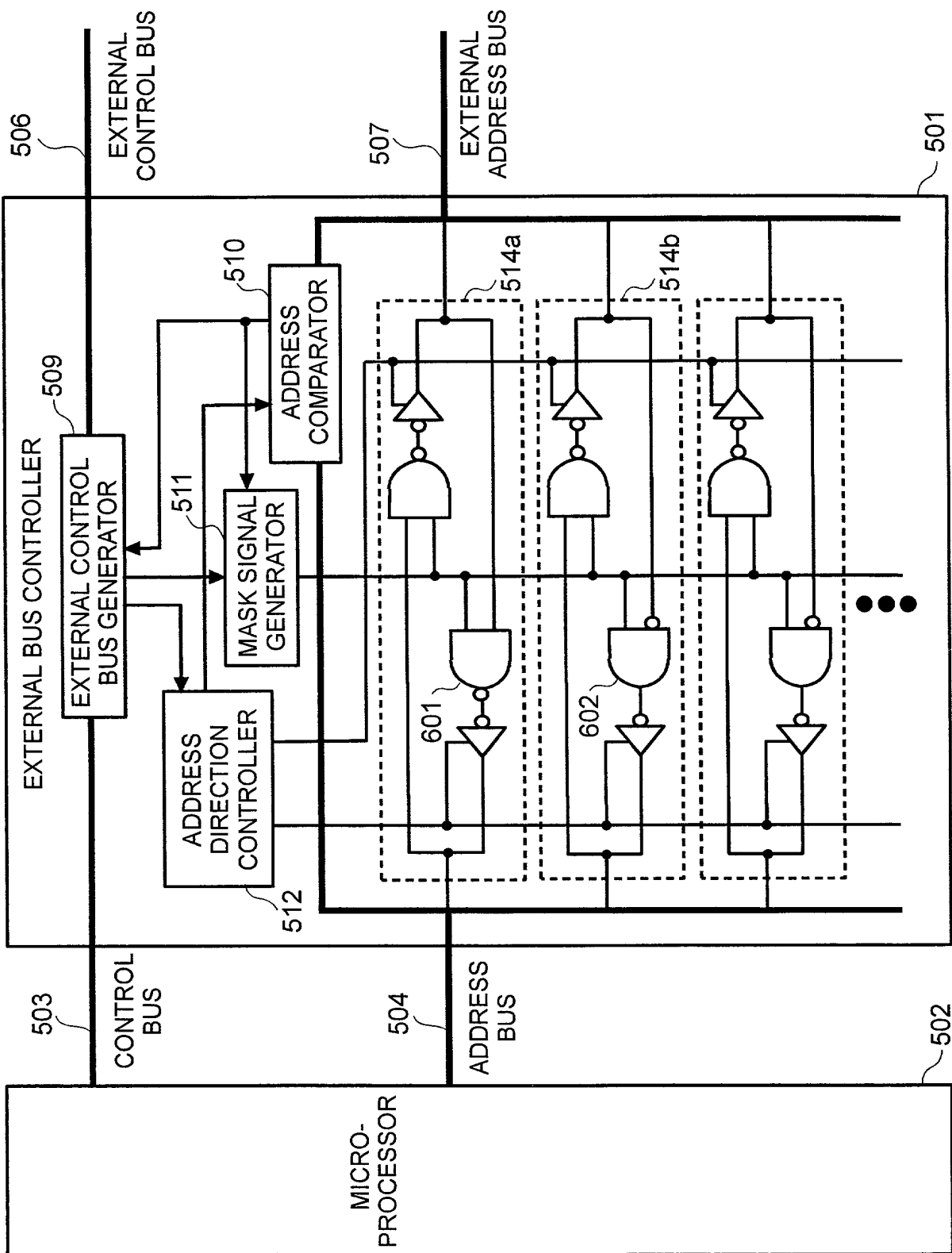


FIG. 6

FIG. 7

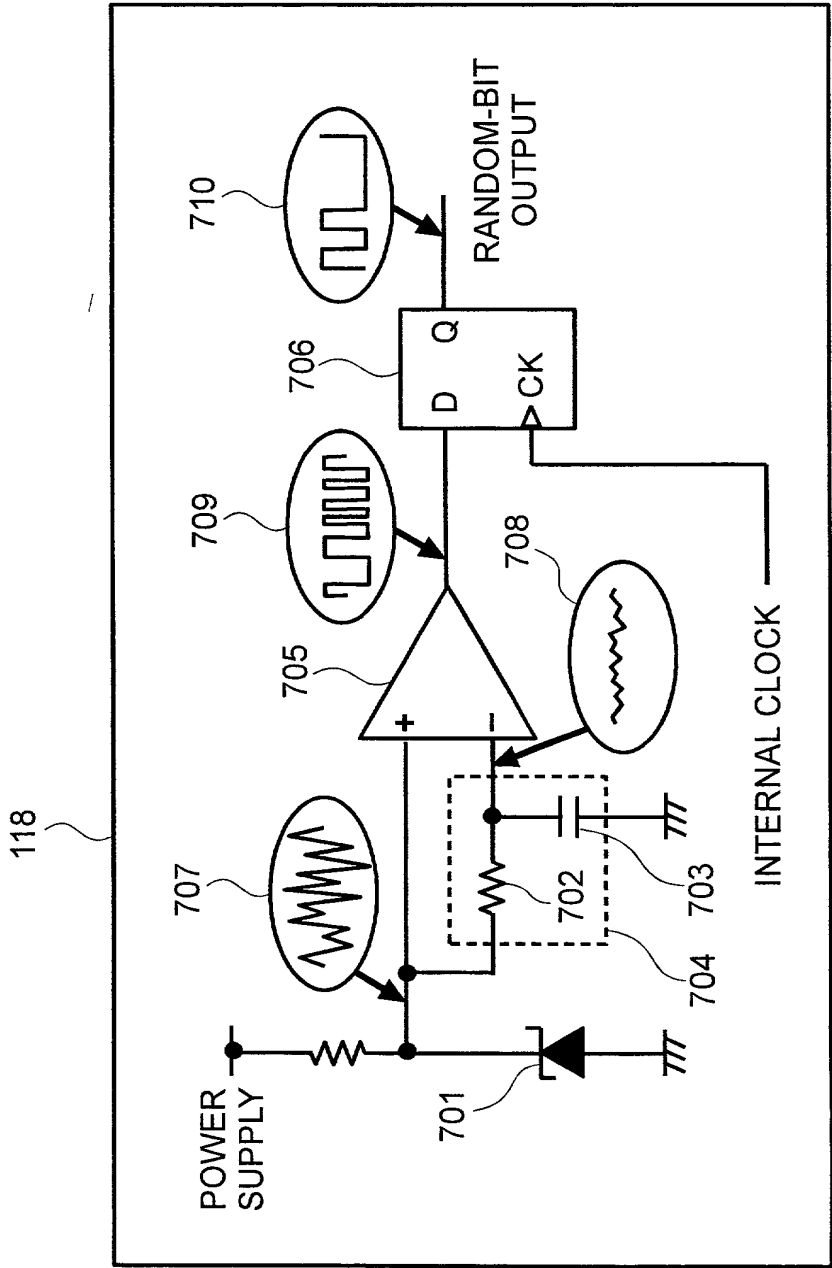


FIG. 8

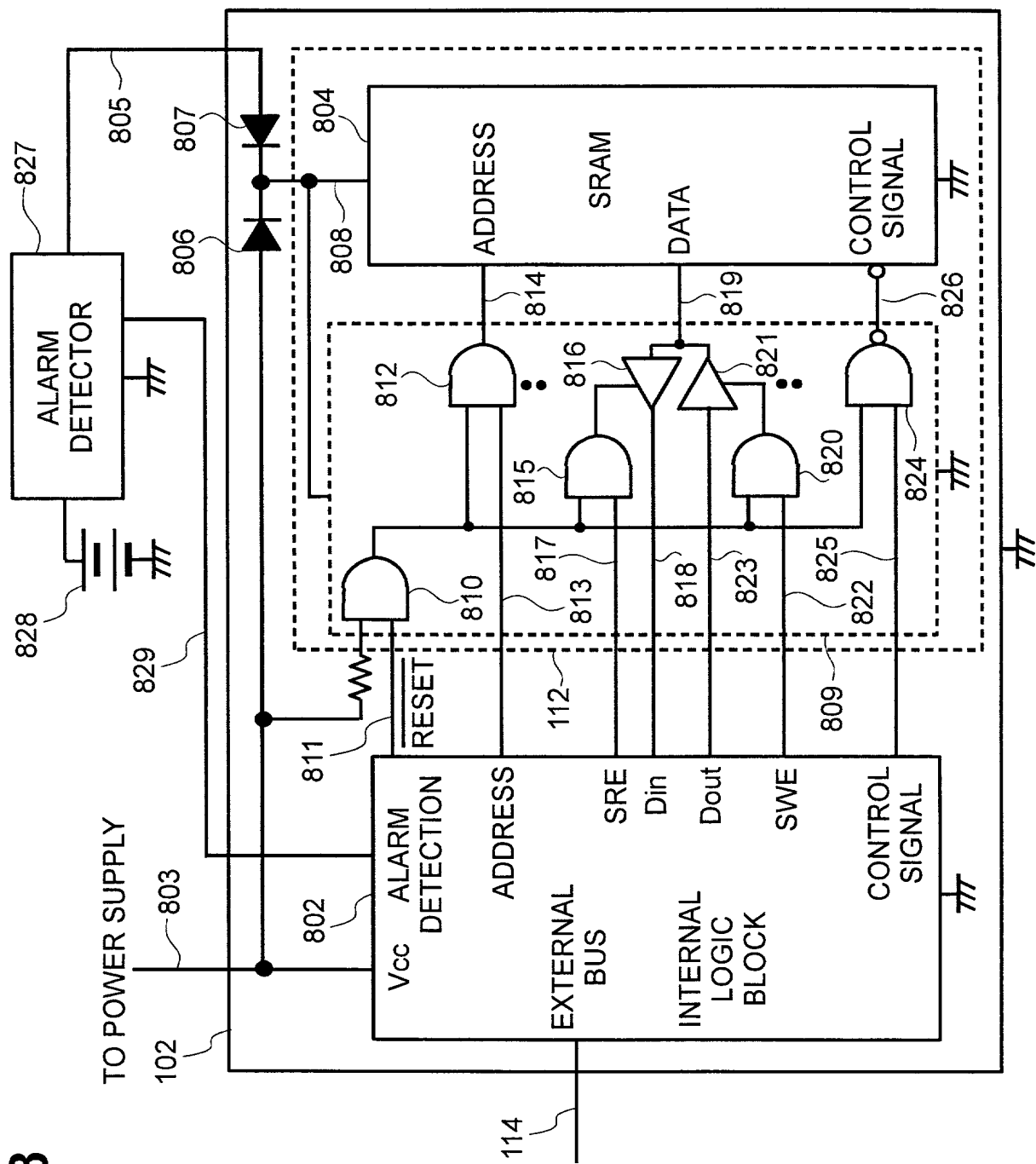
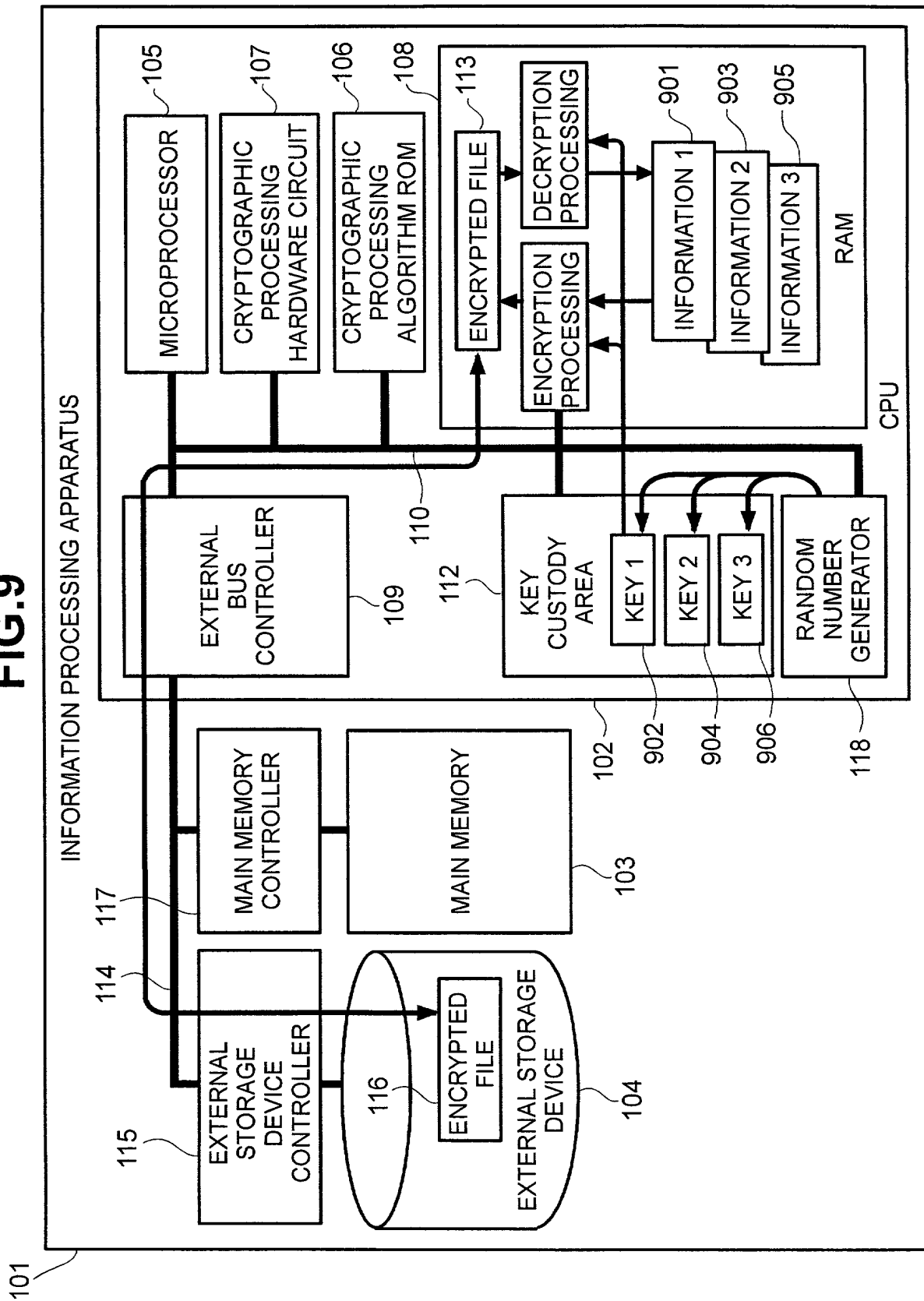
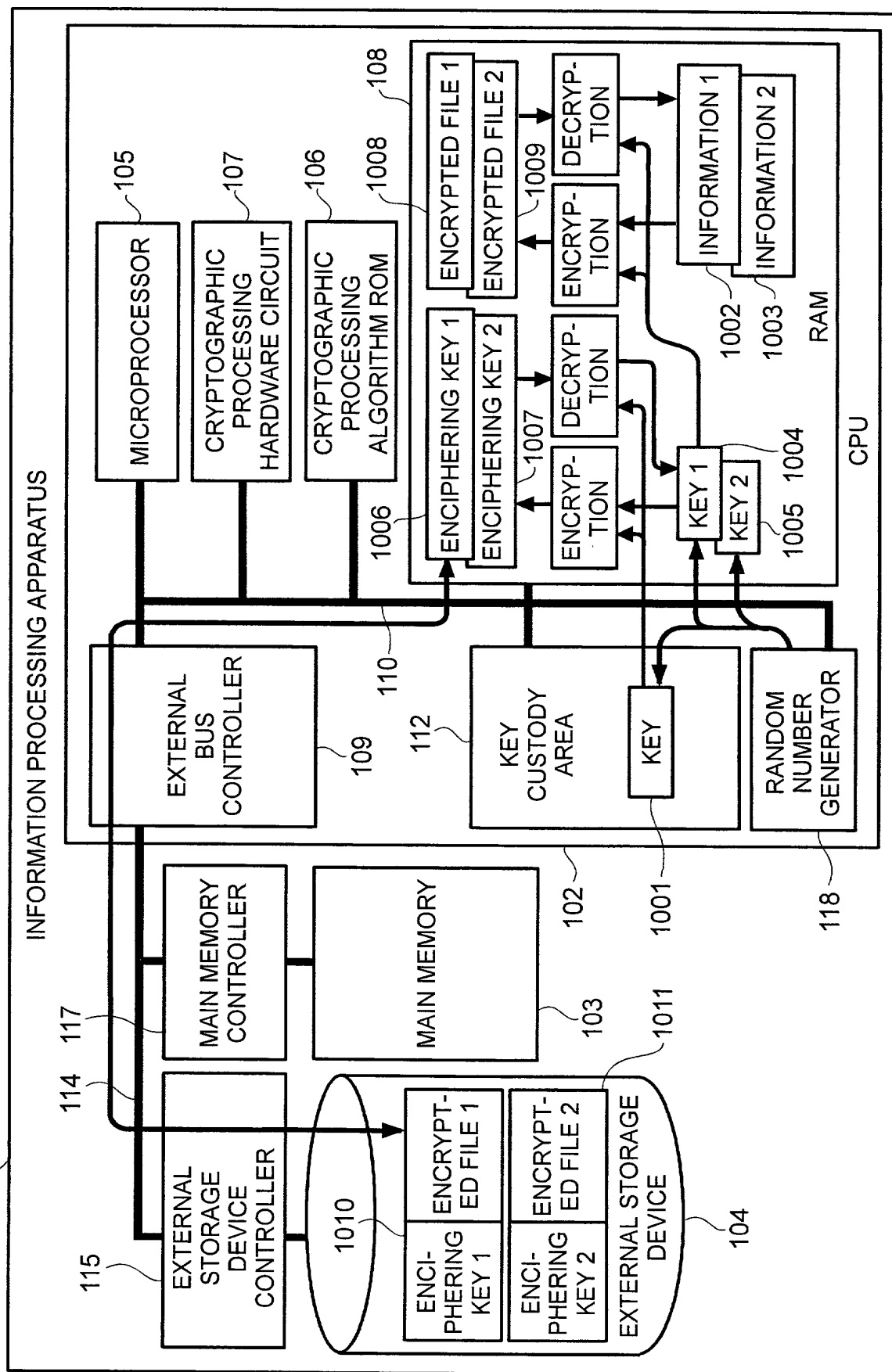


FIG. 9



**FIG. 10**



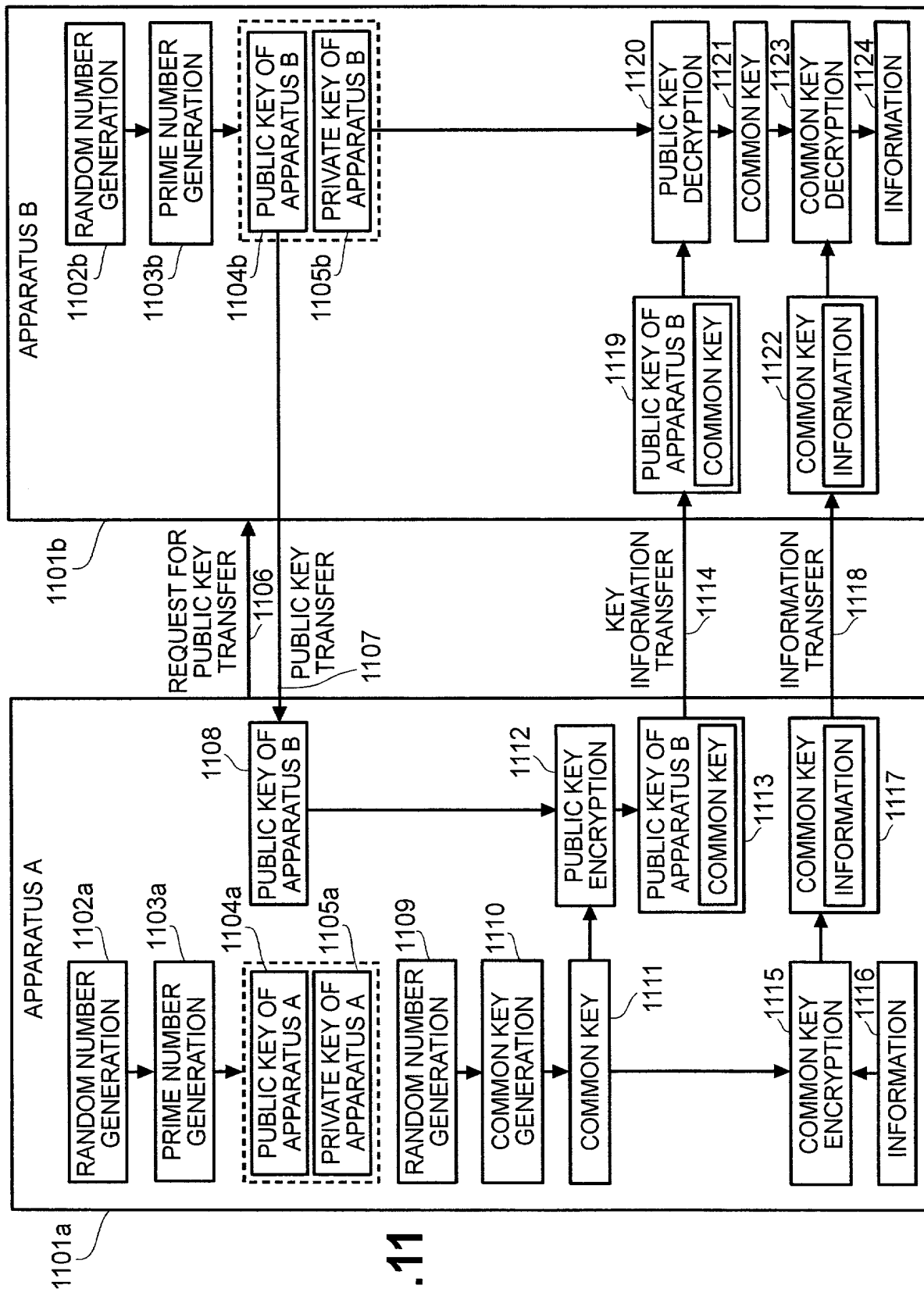
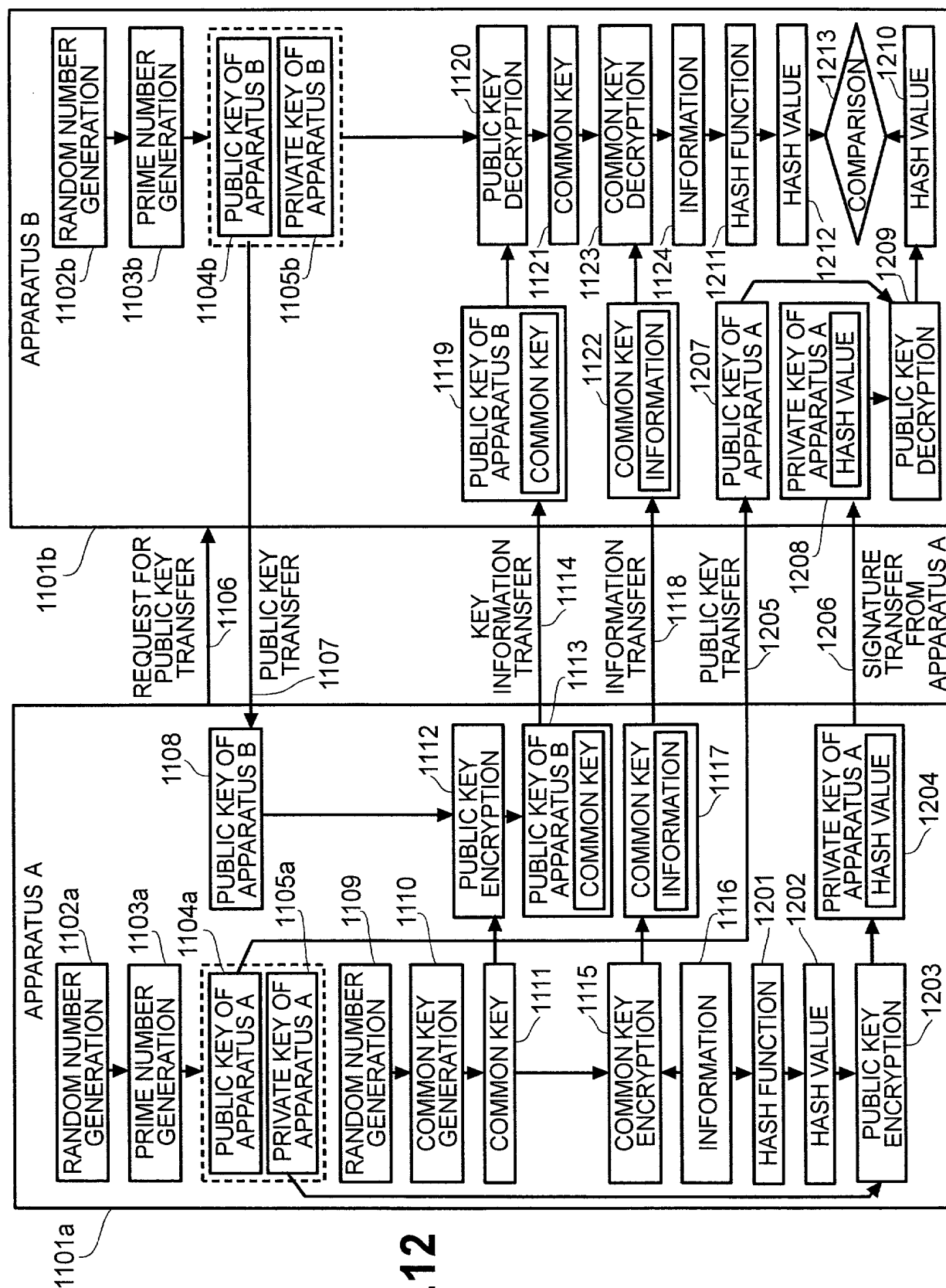


FIG. 11





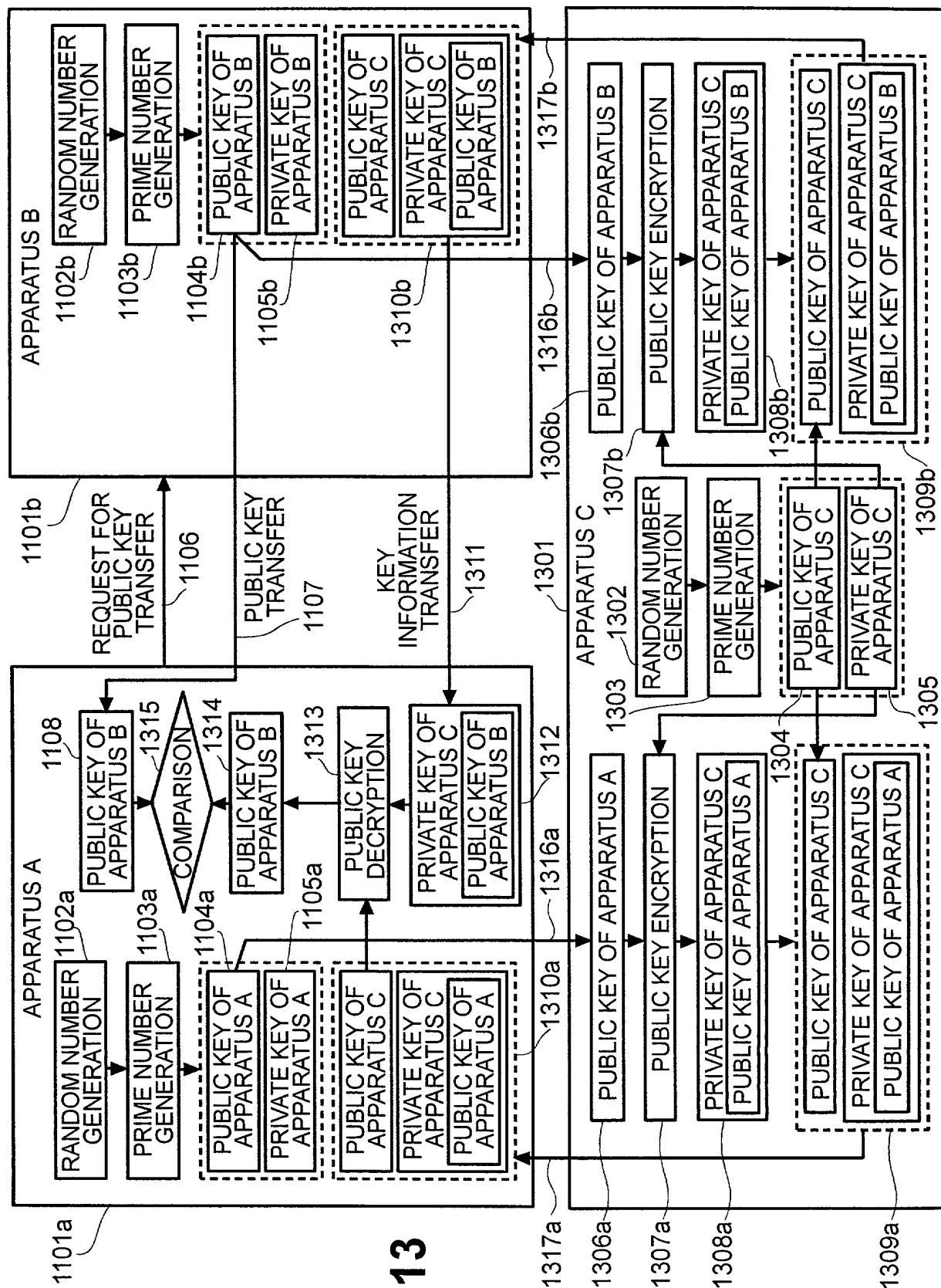


FIG.14

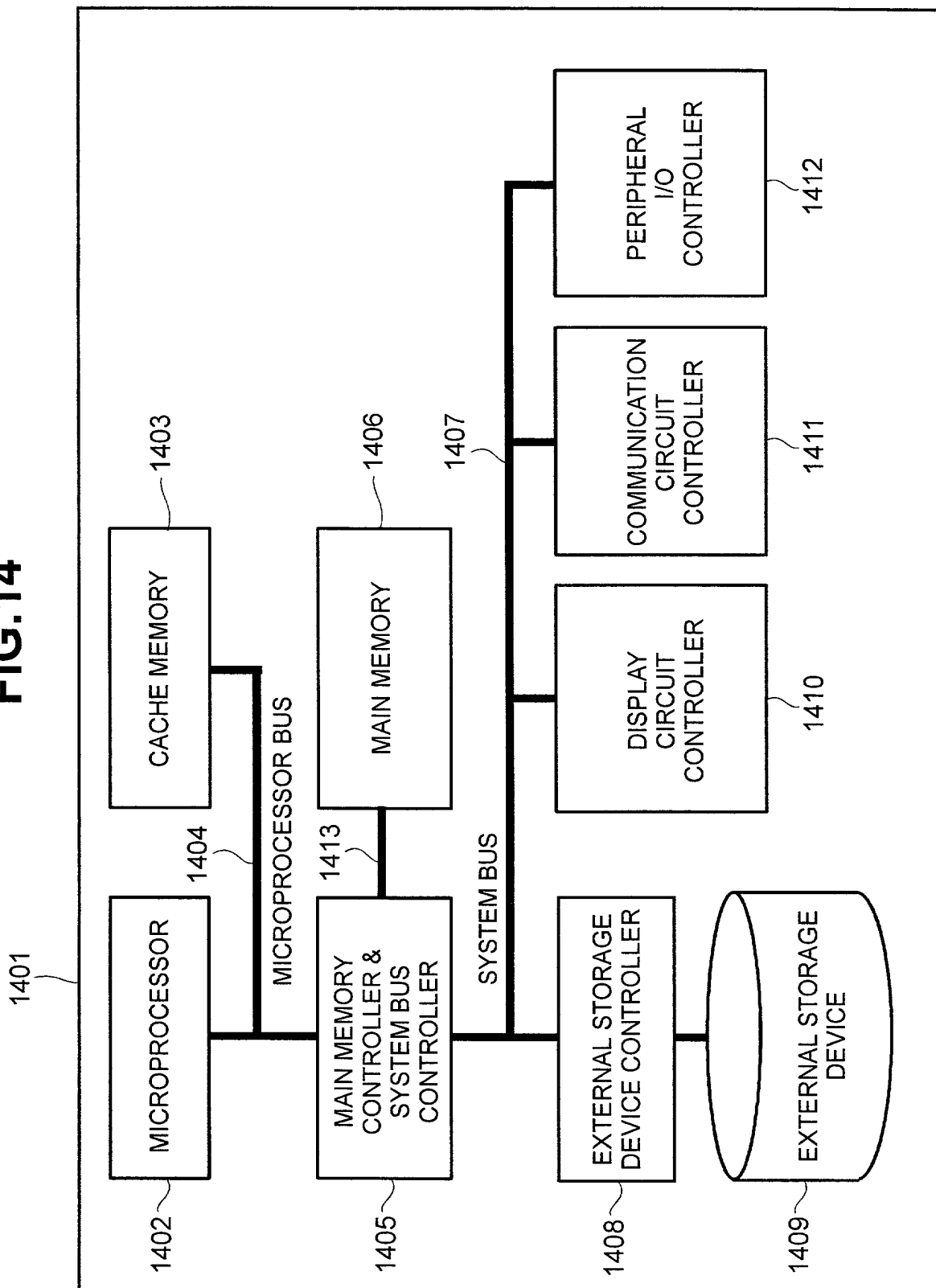


FIG.15

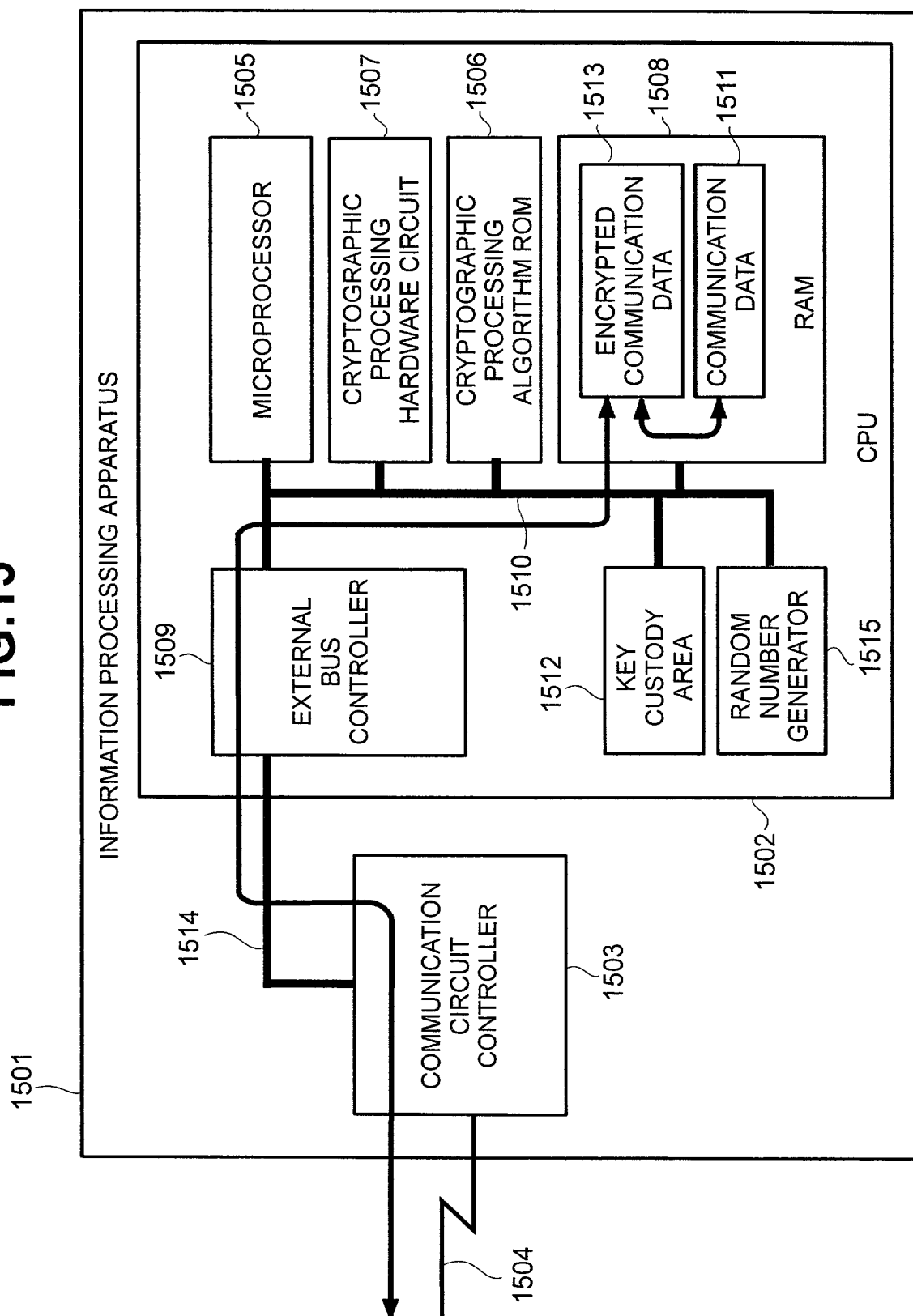
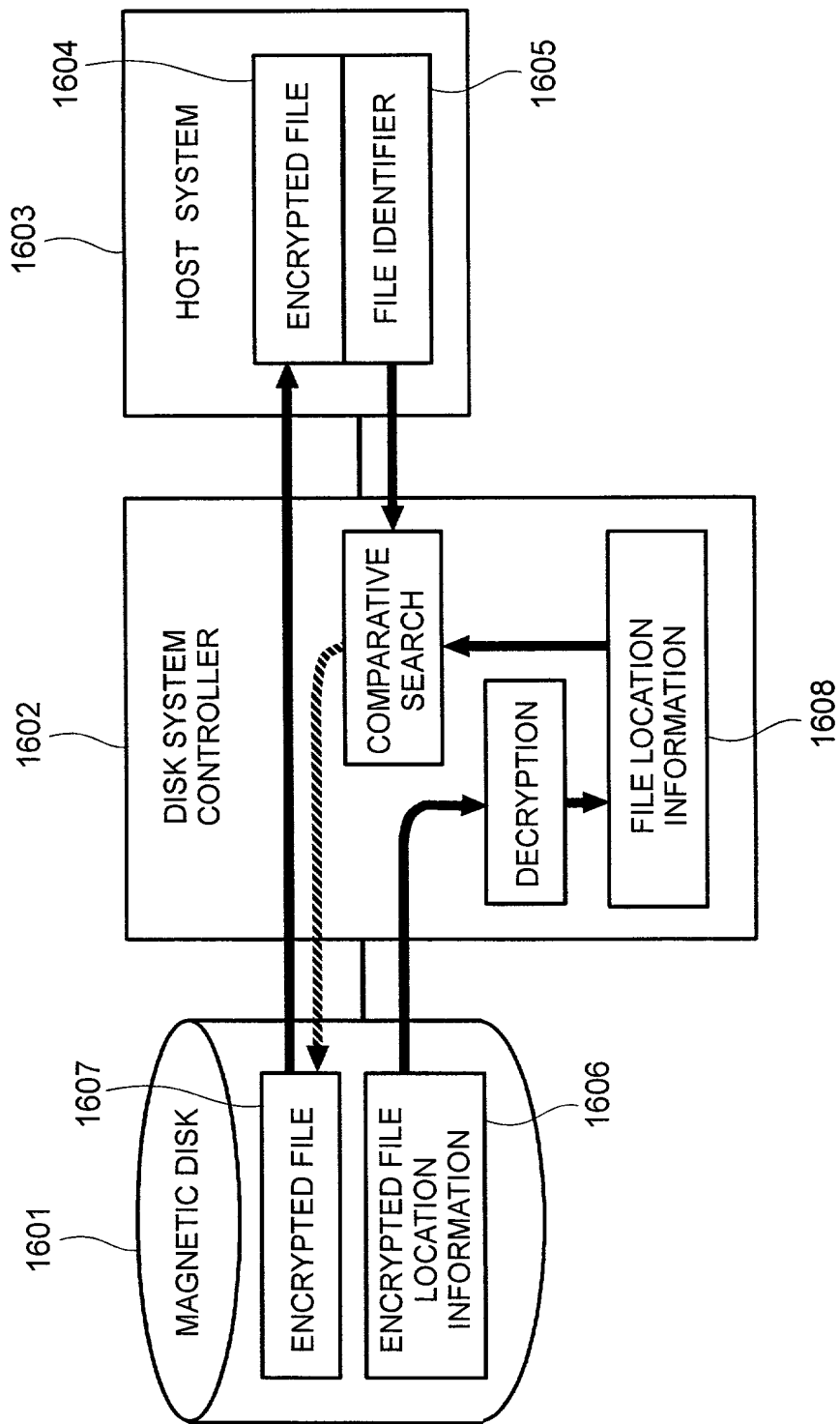


FIG.16



**FIG. 17**

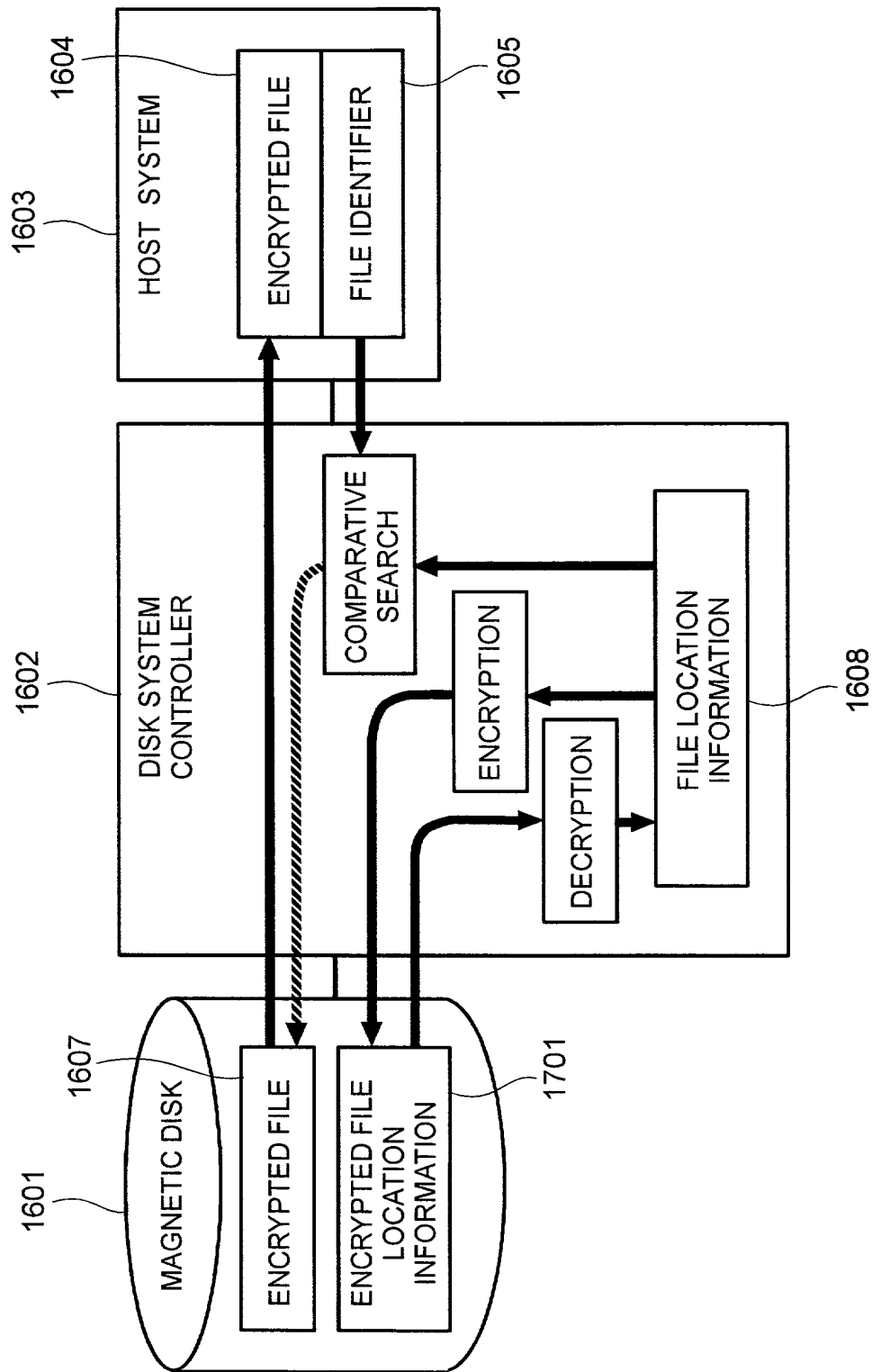
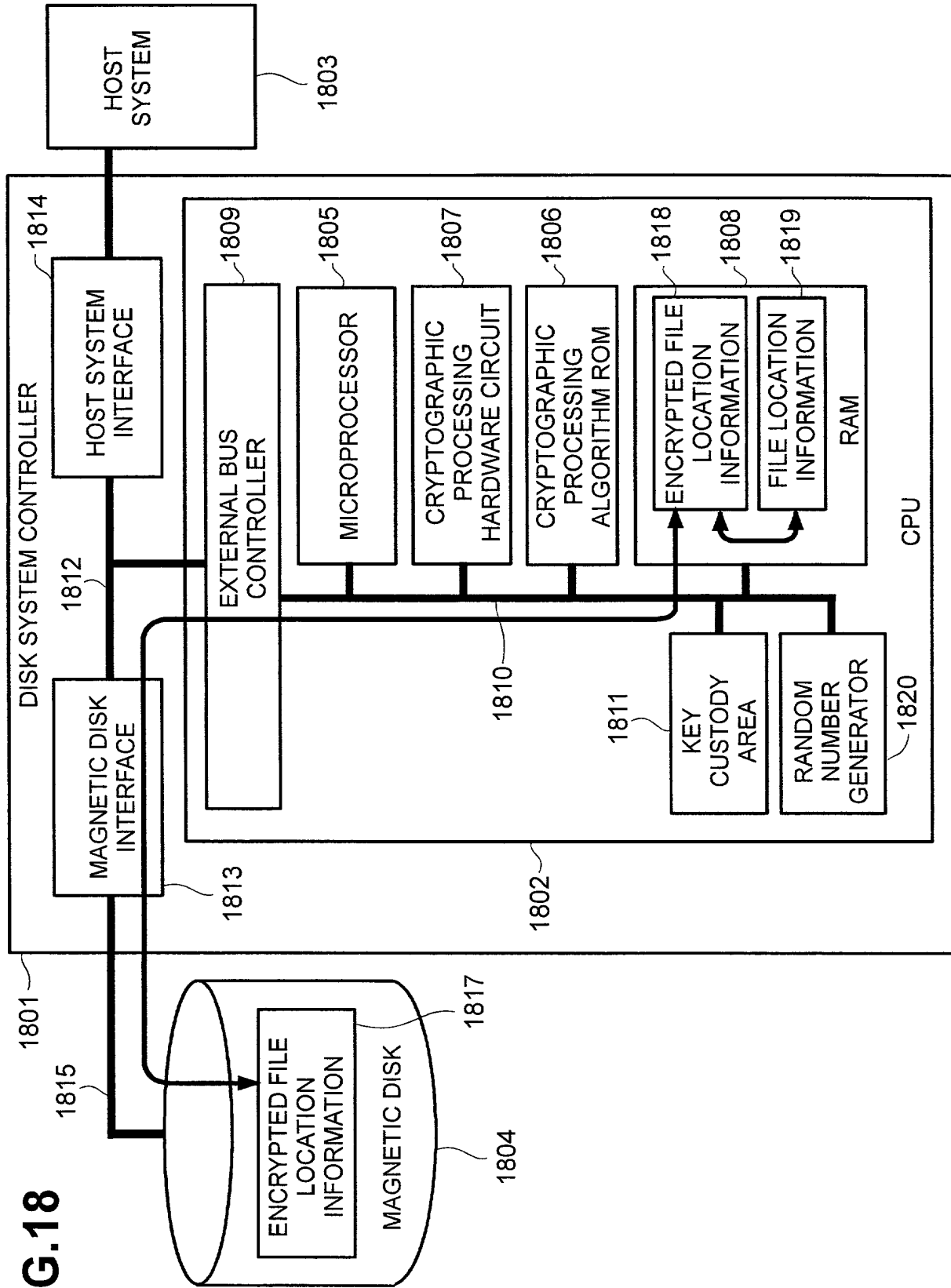


FIG.18



**FIG. 19**

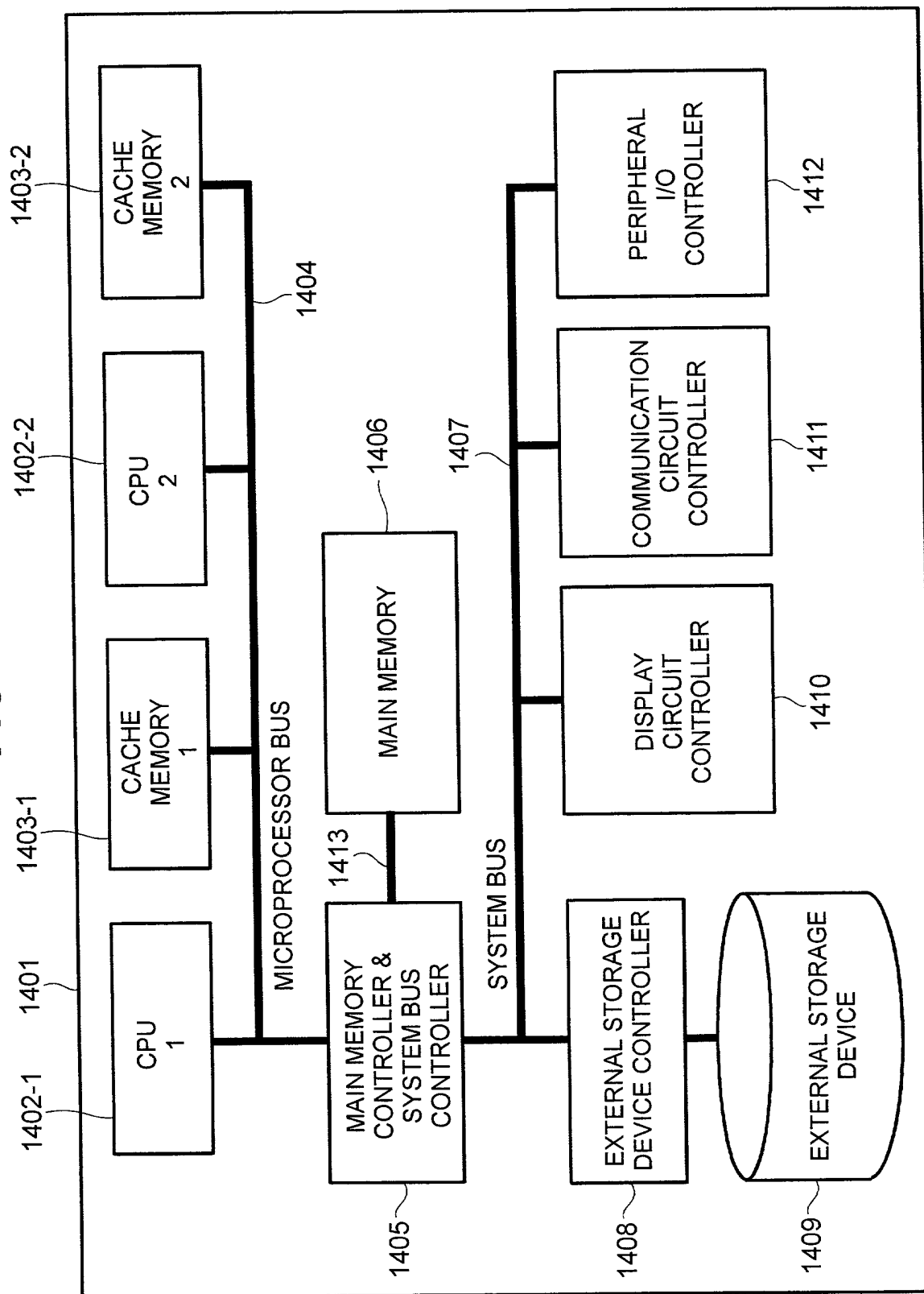
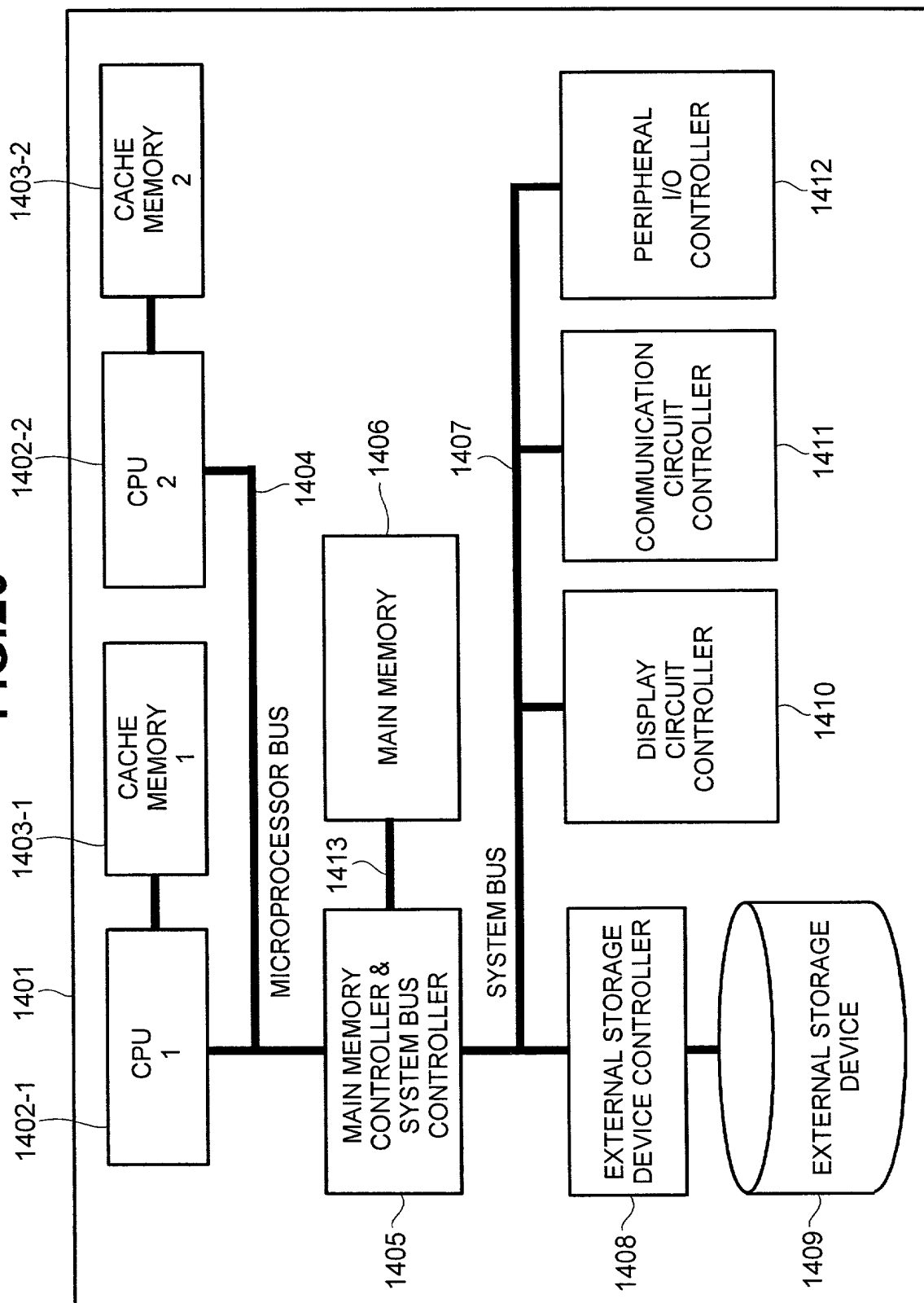


FIG. 20





**FIG. 21**

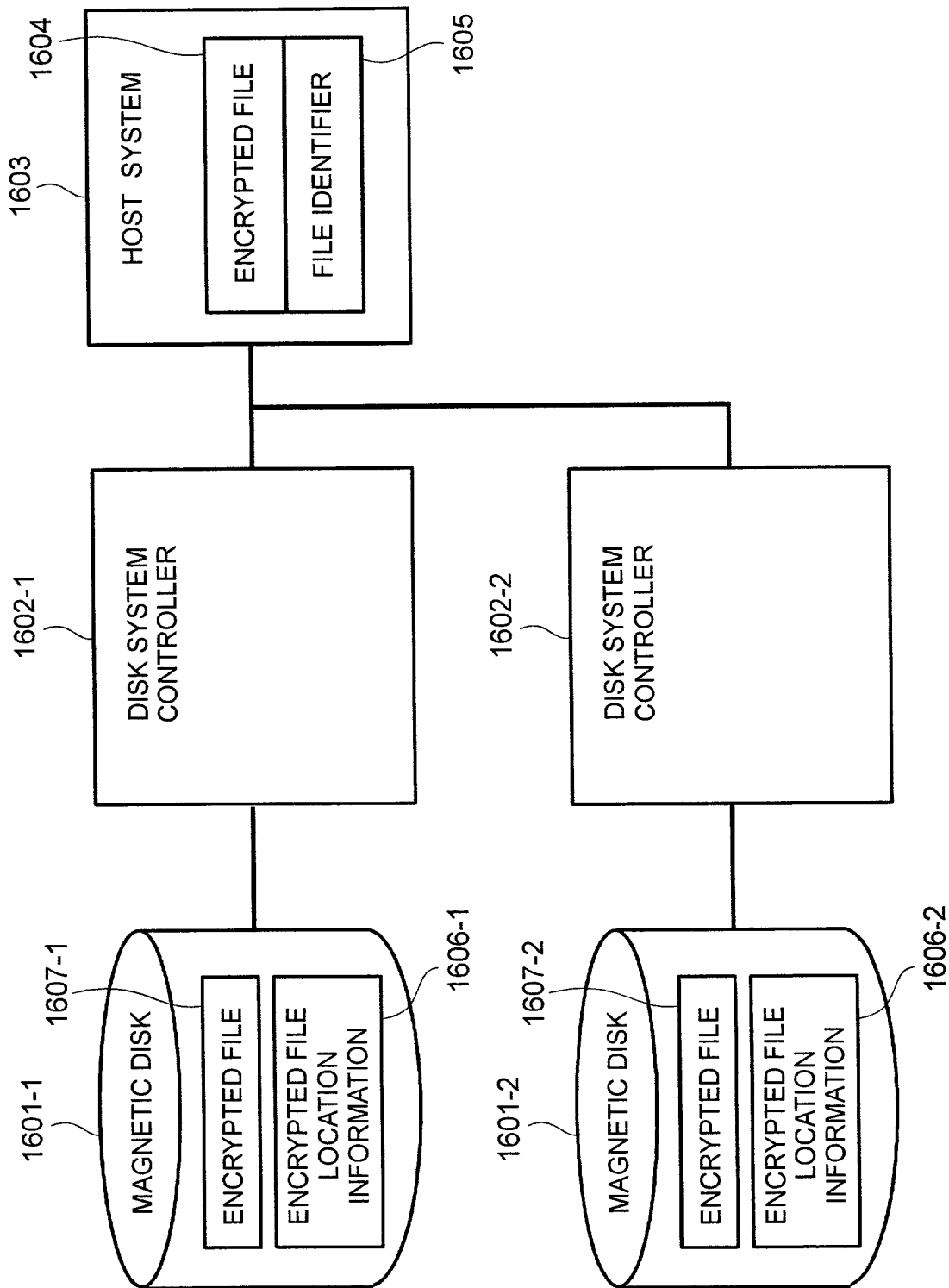


FIG. 22

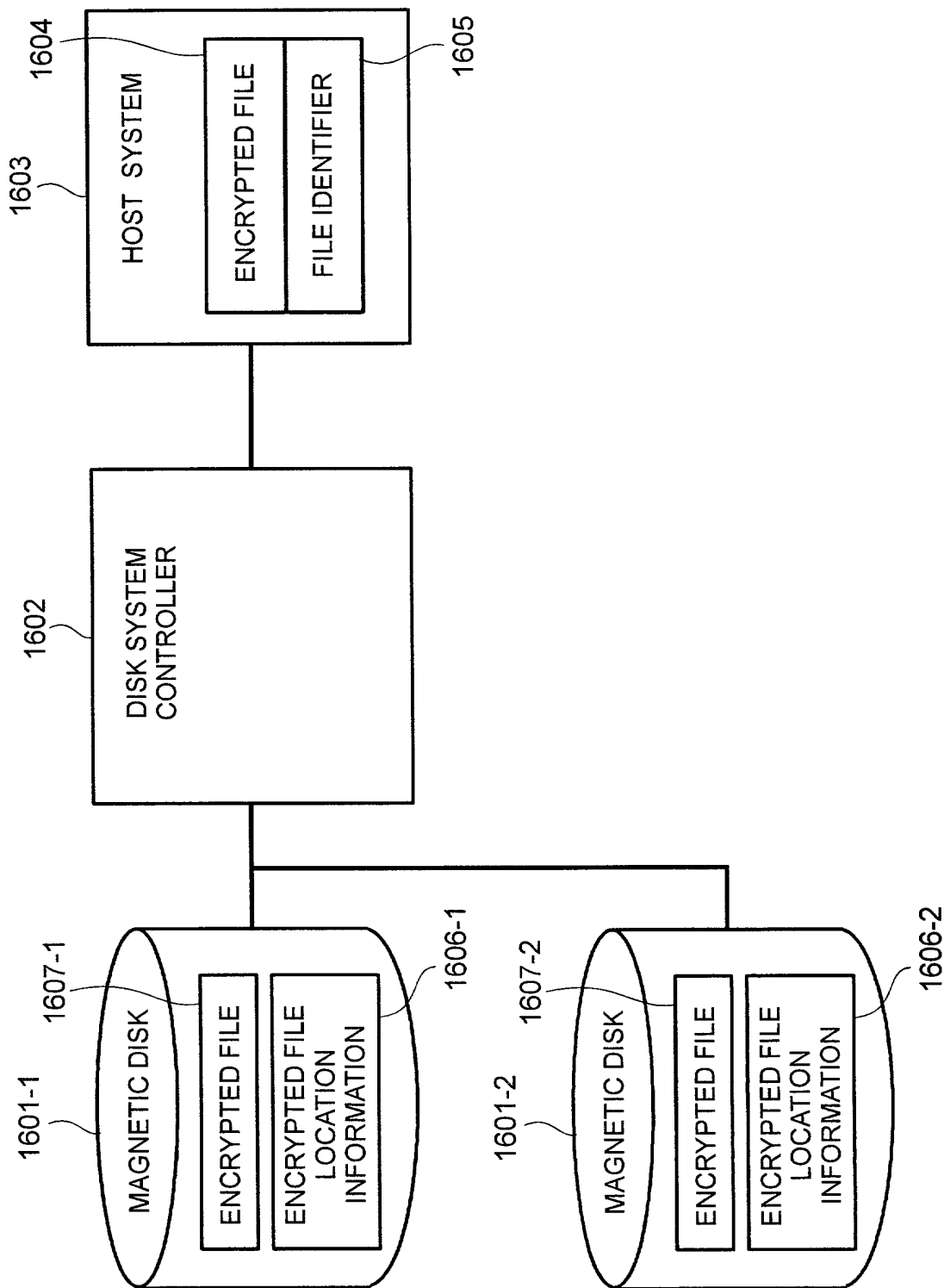
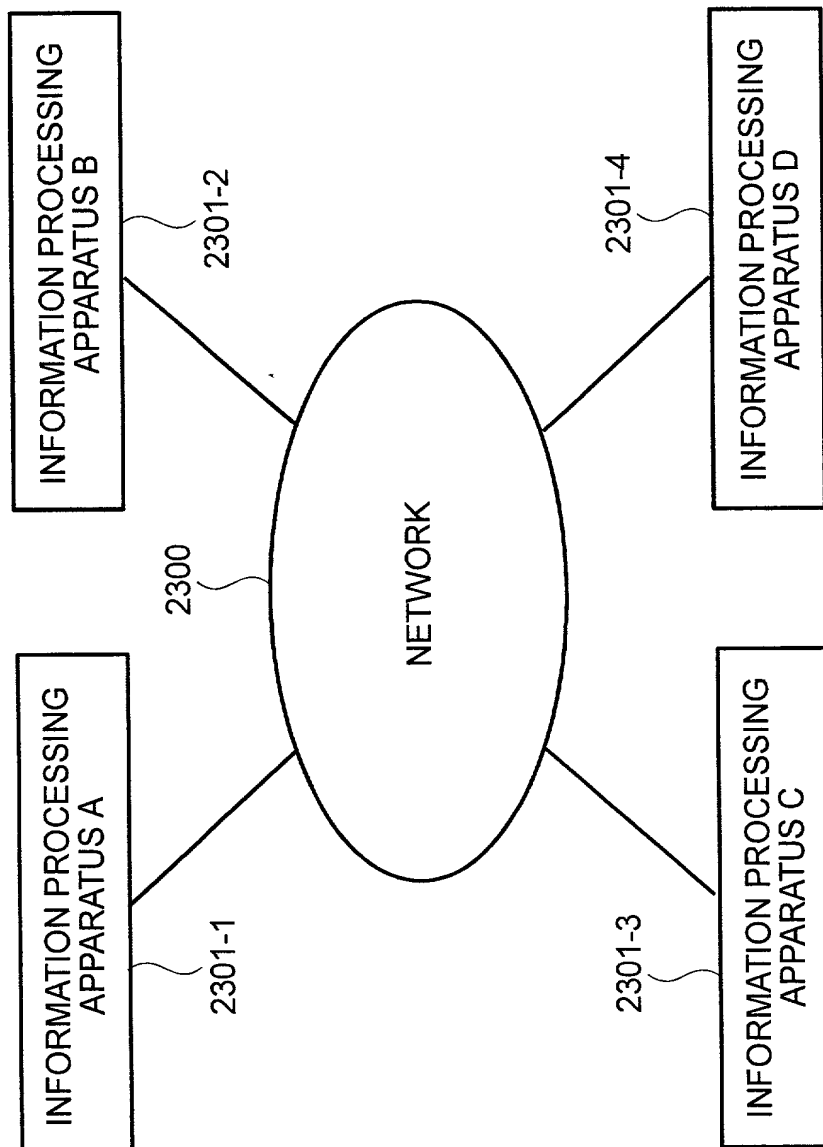


FIG.23



## Declaration and Power of Attorney For Patent Application

## 特許出願宣言書及び委任状

## Japanese Language Declaration

## 日本語宣言書

下記の氏名の発明者として、私は以下の通り宣言します。

As a below named inventor, I hereby declare that:

私の住所、私書箱、国籍は下記の私の氏名の後に記載された通りです。

My residence, post office address and citizenship are as stated next to my name.

下記の名称の発明に関して請求範囲に記載され、特許出願している発明内容について、私が最初かつ唯一の発明者（下記の氏名が一つの場合）もしくは最初かつ共同発明者であると（下記の名称が複数の場合）信じています。

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

## INFORMATION PROCESSING APPARATUS

上記発明の明細書（下記の欄で×印がついていない場合は、本書に添付）は、

The specification of which is attached hereto unless the following box is checked:

☐ \_\_月\_\_日に提出され、米国出願番号または特許協定条約国際出願番号を\_\_\_\_とし、  
(該当する場合) \_\_\_\_\_に訂正されました。

☒ was filed on 6/March/2000  
as United States Application Number or  
PCT International Application Number  
PCT/JP00/01333 and was amended on  
\_\_\_\_\_ (if applicable).

私は、特許請求範囲を含む上記訂正後の明細書を検討し、内容を理解していることをここに表明します。

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment referred to above.

私は、連邦規則法典第37編第1条56項に定義されるとおり、特許資格の有無について重要な情報を開示する義務があることを認めます。

I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56.

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

# Japanese Language Declaration

(日本語宣言書)

私は、米国法典第35編119条(a)-(d)項又は365条(b)項に基き下記の、米国以外の国の少なくとも一カ国を指定している特許協力条約365(a)項に基き国際出願、又は外国での特許出願もしくは発明者証の出願についての外国優先権をここに主張するとともに、優先権を主張している、本出願の前に出願された特許または発明者証の外国出願を以下に、枠内をマークすることで、示している。

## Prior Foreign Application(s)

外国での先行出願

PCT/JP99/01402	PCT
(Number) (番号)	(Country) (国名)
(Number) (番号)	(Country) (国名)

I hereby claim foreign priority under Title 35, United States Code, Section 119 (a)-(d) or 365(b) of any foreign application(s) for patent or inventor's certificate, or 365(a) of any PCT international application which designated at least one country other than the United States, listed below and have also identified below, by checking the box, any foreign application for patent or inventor's certificate, or PCT International application having a filing date before that of the application on which priority is claimed.

Priority Not Claimed  
優先権主張なし

I hereby claim the benefit under Title 35, United States Code, Section 119(e) of any United States provisional application(s) listed below.

私は、第35編米国法典119条(e)項に基いて下記の米国特許出願規定に記載された権利をここに主張いたします。

(Application No.) (出願番号)	(Filing Date) (出願日)
-----------------------------	------------------------

(Application No.) (出願番号)	(Filing Date) (出願日)
-----------------------------	------------------------

私は、下記の米国法典第35編120条に基いて下記の米国特許出願に記載された権利、又は米国を指定している特許協力条約365条(c)に基き権利をここに主張します。また、本出願の各請求範囲の内容が米国法典第35編112条第1項又は特許協力条約で規定された方法で先行する米国特許出願に開示されていない限り、その先行米国出願書提出日以降で本出願書の日本国内または特許協力条約国際提出日までの期間中に入手された、連邦規則法典第37編1条56項で定義された特許資格の有無に関する重要な情報について開示義務があることを認識しています。

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s), or 365(c) of any PCT international application designating the United States, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of Title 35, United States Code Section 112, I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of application.

(Application No.) (出願番号)	(Filing Date) (出願日)
-----------------------------	------------------------

(Status: Patented, Pending, Abandoned) (現況: 特許許可済、係属中、放棄済)
---

(Application No.) (出願番号)	(Filing Date) (出願日)
-----------------------------	------------------------

(Status: Patented, Pending, Abandoned) (現況: 特許許可済、係属中、放棄済)
---

私は、私自身の知識に基づいて本宣言書中で私が行なう表明が真実であり、かつ私の入手した情報と私の信じるところに基き、かつ表明が全て真実であると信じていること、さらに故意になされた虚偽の表明及びそれと同等の行為は米国法典第18編第1001条に基づき、罰金または拘禁、もしくはその両方により処罰されること、そしてそのような故意による虚偽の声明を行えば、出願した、又は既に許可された特許の有効性が失われることを認識し、よってここに上記のごとく宣誓を致します。

I hereby declare that all statements made herein of my own

knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

### Japanese Language Declaration (日本語宣言書)

委任状： 私は下記の発明者として、本出願に関する一切の手続きを米特許商標局に対して遂行する弁理士または代理人として、下記の者を指名いたします。(弁理士、または代理人の氏名及び登録番号を明記のこと)

POWER OF ATTORNEY: As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith (list name and registration number)

John R. Mattingly, Reg.No.30,293; Daniel J. Stanger, Reg. No.32,846; Shrinath Malur, Reg. No.34,663; Gene W. Stockman, Reg. No.21,021; Jeffrey M.Ketchum, Reg. No.31,174; and Scott W. Brickner, Reg. No.34,553.

書類送付先

Send Correspondence to:

Mattingly, Stanger & Malur, P.C.

104 East Hume Avenue

Alexandria, Virginia 22301

直接電話連絡先： (氏名及び電話番号)

Direct Telephone Calls to: (name and telephone number)

Telephone: (703) 684-1120

Fax: (703) 684-1157

唯一または第一発明者

Full name of sole or first inventor

Jun KITAHARA

発明者の署名

日付

Inventor's signature

Date Jun. 6, 2001

Jun Kitahara

住所

Residence

Kawasaki, Japan

国籍

Citizenship

Japan

私書箱

Post Office Address

c/o Hitachi, Ltd., Intellectual Property Group

New Marunouchi Bldg. 5-1, Marunouchi 1-chome,

Chiyoda-ku, Tokyo 100-8220, Japan

(第二以降の共同発明者についても同様に記載し、署名をすること)

(Supply similar information and signature for second and subsequent joint inventors.)

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

第二共同発明者	2-00	Full name of second joint inventor, if any Takeshi ASAHI
第二共同発明者の署名	日付	Second inventor's signature Date Jun. 6, 2001 Takeshi Asahi
住所		Residence Kawasaki, Japan SPX
国籍		Citizenship Japan
私書箱		Post Office Address c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan
第三共同発明者	3-00	Full name of third joint inventor, if any Toru OWADA
第三共同発明者の署名	日付	Third inventor's signature Date Jun. 6, 2001 Toru Owada
住所		Residence Kawasaki, Japan JPY
国籍		Citizenship Japan
私書箱		Post Office Address c/o Hitachi, Ltd., Intellectual Property Group New Marunouchi Bldg. 5-1, Marunouchi 1-chome, Chiyoda-ku, Tokyo 100-8220, Japan
第四共同発明者		Full name of fourth joint inventor, if any
第四共同発明者の署名	日付	Fourth inventor's signature Date
住所		Residence
国籍		Citizenship
私書箱		Post Office Address
第五共同発明者		Full name of fifth joint inventor, if any
第五共同発明者の署名	日付	Fifth inventor's signature Date
住所		Residence
国籍		Citizenship
私書箱		Post Office Address

(第六以降の共同発明者についても同様に記載し、署名をすること)

(Supply similar

information and signature for sixth and subsequent joint inventors.)